

On the Practical Computational Power of Finite Precision RNNs for Language Recognition

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GRU < LSTM (!?)



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Current State

- RNNs are everywhere
- We don't know too much about the differences between them:
 - Gated RNNs are shown to train better, beyond that:
 - "RNNs are Turing Complete"?

Turing Complete?

On the Computational Power of Neural Nets*

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Turing Complete?

1993 Proof:

1. Requires Infinite Precision:

Uses stack(s), maintained in certain dimension(s) Zeros are pushed using division (using g = g/4 + 1/4) In 32 bits, this reaches the limit after **15** pushes

2. Requires Infinite Time:

Allows processing steps beyond reading input (Not the standard use case!)

unreasonable assumptions!

Turing Complete?



unreasonable assumptions!

What happens on real hardware and real use-cases?

Real Use

- Gated architectures have the best performance
 - LSTM and GRU are most popular
 - Of these, the choice between them is unclear

Main Result

We accept all RNN types can simulate DFAs

We show that LSTMs and IRNNs can also count

And that the GRU and SRNN cannot

Power of Counting

Practical

In NMT: LSTM better at capturing target length

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Theoretical

Finite State Machines vs Counter Machines

K-Counter Machines (SKCMs)

Fischer, Meyer, Rosenberg - 1968

- Similar to finite automata, but also maintain k counters
- A counter has 4 operations: inc/dec by one, do nothing, reset
- Counters are observed by comparison to zero



Counting Machines and Chomsky Hierarchy











SKCMs cross the Chomsky Hierarchy!



Summary so Far

- Counters give additional formal power
- We claimed that LSTM can count and GRU cannot

Summary so Far

- Counters give additional formal power
- We claimed that LSTM can count and GRU cannot
- Let's see why

$$z_t = \sigma(W^z x_t + U^z h_{t-1} + b^z)$$

$$r_t = \sigma(W^r x_t + U^r h_{t-1} + b^r)$$

$$\tilde{h}_t = \tanh(W^h x_t + U^h (r_t \circ h_{t-1}) + b^h)$$

$$h_t = z_t \circ h_{t-1} + (1 - z_t) \circ \tilde{h}_t$$

$$\begin{split} f_t &= \sigma(W^f x_t + U^f h_{t-1} + b^f) \\ i_t &= \sigma(W^i x_t + U^i h_{t-1} + b^i) \\ o_t &= \sigma(W^o x_t + U^o h_{t-1} + b^o) \\ \tilde{c}_t &= \tanh(W^c x_t + U^c h_{t-1} + b^c) \\ c_t &= f_t \circ c_{t-1} + i_t \circ \tilde{c}_t \\ h_t &= o_t \circ g(c_t) \end{split}$$

$$\begin{aligned} z_t &= \sigma(W^z x_t + U^z h_{t-1} + b^z) & \text{gates} \\ r_t &= \sigma(W^r x_t + U^r h_{t-1} + b^r) \\ \tilde{h}_t &= \tanh(W^h x_t + U^h (r_t \circ h_{t-1}) + b^h) \\ \tilde{h}_t &= tanh(W^h x_t + U^h (r_t \circ h_{t-1}) + b^h) \\ h_t &= z_t \circ h_{t-1} + (1 - z_t) \circ \tilde{h}_t \\ & \text{candidate} \\ & \text{vectors} \end{aligned} \qquad \begin{aligned} f_t &= \sigma(W^j x_t + U^j h_{t-1} + b^j) \\ o_t &= \sigma(W^o x_t + U^o h_{t-1} + b^o) \\ \tilde{c}_t &= tanh(W^c x_t + U^c h_{t-1} + b^c) \\ c_t &= f_t \circ c_{t-1} + i_t \circ \tilde{c}_t \\ h_t &= o_t \circ g(c_t) \end{aligned}$$



→ gates → $f_t \in (0,1)$ $z_t \in (0,1)$ $i_t \in (0,1)$ $r_t \in (0,1)$ $o_t \in (0,1)$ $\tilde{h}_t \in (-1,1)$ $\tilde{c}_t \in (-1,1)$ $h_t = z_t \circ h_{t-1} + (1 - z_t) \circ h_t$ $c_t = f_t \circ c_{t-1} + i_t \circ \tilde{c}_t$ candidate $h_t = o_t \circ g(c_t)$ vectors update functions

 $z_t \in (0,1)$ $r_t \in (0,1)$ $\tilde{h}_t \in (-1,1)$ $h_t = z_t \circ h_{t-1} + (1-z) \circ \tilde{h}_t$ $f_{t} \in (0,1)$ $i_{t} \in (0,1)$ $o_{t} \in (0,1)$ $\tilde{c}_{t} \in (-1,1)$ $c_{t} = f_{t} \circ c_{t-1} + i_{t} \circ \tilde{c}_{t}$ $h_{t} = o_{t} \circ g(c_{t})$

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Interpolation

 $z_t \in (0,1)$ $r_t \in \textbf{Bounded!}$ $\tilde{h}_t \in (-1,1)$ $h_t = z_t \circ h_{t-1} + (1-z) \circ \tilde{h}_t$

Interpolation

 $f_{t} \in (0,1)$ $i_{t} \in (0,1)$ $o_{t} \in (0,1)$ $\tilde{c}_{t} \in (-1,1)$ $c_{t} = f_{t} \circ c_{t-1} + i_{t} \circ \tilde{c}_{t}$ $h_{t} = o_{t} \circ g(c_{t})$

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Interpolation

 $f_{t} \in (0,1)$ $i_{t} \in (0,1)$ $o_{t} \in (0,1)$ $\tilde{c}_{t} \in (-1,1)$ $c_{t} = f_{t} \circ c_{t-1} + i_{t} \circ \tilde{c}_{t}$ $h_{t} = o_{t} \circ g(c_{t})$

Addition

 $z_{t} \in (0,1)$ $r_{t} \in \textbf{Bounded!}$ $\tilde{h}_{t} \in (-1,1)$ $h_{t} = z_{t} \circ h_{t-1} + (1-z) \circ \tilde{h}_{t}$

Interpolation

 $f_t \approx 1$ $i_t \approx 1$ $o_t \in (0,1)$ $\tilde{c}_t \in (-1,1)$ $c_t \approx c_{t-1} + \tilde{c}_t$ $h_t = o_t \circ g(c_t)$

Addition

$$c_t = f_t \circ c_{t-1} + i_t \circ \tilde{c}_t$$

 $z_{t} \in (0,1)$ $r_{t} \in \textbf{Bounded!}$ $\tilde{h}_{t} \in (-1,1)$ $h_{t} = z_{t} \circ h_{t-1} + (1-z) \circ \tilde{h}_{t}$

Interpolation

 $f_t \approx 1$ $i_t \approx 1$ $o_t \in (0,1)$ $\tilde{c}_t \approx 1$ $c_t \approx c_{t-1} + 1$ $h_t = o_t \circ g(c_t)$

Increase by 1

$$c_t = f_t \circ c_{t-1} + i_t \circ \tilde{c}_t$$

 $z_{t} \in (0,1)$ $r_{t} \in \textbf{Bounded!}$ $\tilde{h}_{t} \in (-1,1)$ $h_{t} = z_{t} \circ h_{t-1} + (1-z) \circ \tilde{h}_{t}$

Interpolation

 $f_t \approx 1$ $i_t \approx 1$ $o_t \in (0,1)$ $\tilde{c}_t \approx -1$ $c_t \approx c_{t-1} - 1$ $h_t = o_t \circ g(c_t)$

Decrease by 1

$$c_t = f_t \circ c_{t-1} + i_t \circ \tilde{c}_t$$

 $z_{t} \in (0,1)$ $r_{t} \in \textbf{Bounded!}$ $\tilde{h}_{t} \in (-1,1)$ $h_{t} = z_{t} \circ h_{t-1} + (1-z) \circ \tilde{h}_{t}$

Interpolation

 $f_t \approx 1$ $i_t \approx 0$ $o_t \in (0,1)$ $\tilde{c}_t \in (-1,1)$ $c_t \approx c_{t-1}$ $h_t = o_t \circ g(c_t)$

Do Nothing

$$c_t = f_t \circ c_{t-1} + i_t \circ \tilde{c}_t$$

 $z_{t} \in (0,1)$ $r_{t} \in \textbf{Bounded!}$ $\tilde{h}_{t} \in (-1,1)$ $h_{t} = z_{t} \circ h_{t-1} + (1-z) \circ \tilde{h}_{t}$

Interpolation

 $f_t \approx 0$ $i_t \approx 0$ $o_t \in (0,1)$ $\tilde{c}_t \in (-1,1)$ $c_t \approx 0$ $h_t = o_t \circ g(c_t)$

Reset

$$c_t = f_t \circ c_{t-1} + i_t \circ \tilde{c}_t$$

 $z_{t} \in (0,1)$ $r_{t} \in \textbf{Bounded!}$ $\tilde{h}_{t} \in (-1,1)$ $h_{t} = z_{t} \circ h_{t-1} + (1-z) \circ \tilde{h}_{t}$

Interpolation

 $f_t \approx 0$ $i_t \approx 0$ $o_t \in \textbf{Can Count!}$ $o_t \in (-1,1)$ $c_t \approx 0$ $h_t = o_t \circ g(c_t)$

Reset

$$c_t = f_t \circ c_{t-1} + i_t \circ \tilde{c}_t$$

SRNN

IRNN

 $h_{t} = \sigma_{h}(W_{h}x_{t} + U_{h}h_{t-1} + b_{h}) \qquad h_{t} = \max(0, W_{h}x_{t} + U_{h}h_{t-1} + b_{h})$

SRNN

IRNN

 $h_t = \sigma_h(W_h x_t + U_h h_{t-1} + b_h) \in (0,1) \qquad h_t = \max(0, W_h x_t + U_h h_{t-1} + b_h)$



SRNN

IRNN



(subtraction in parallel, also increasing, counter)

SRNN

IRNN



(subtraction in parallel, also increasing, counter)

So:

- LSTM can count!
- GRU cannot
- Counting gives greater computational power

Trained $a^n b^n$, (on positive examples up to length 100)

Activations on $a^{1000}b^{1000}$:



GRU

Trained $a^n b^n$, (on positive examples up to length 100)

Activations on $a^{1000}b^{1000}$:



• Took much longer to train

Trained $a^n b^n$, (on positive examples up to length 100)

Activations on $a^{1000}b^{1000}$:



GRU

GRU:

- Took much longer to train
- Did not generalise even within training domain
 - begin failing at n=39 (vs 257 for LSTM)

Trained $a^n b^n$, (on positive examples up to length 100)

Activations on $a^{1000}b^{1000}$:



GRU:

- Took much longer to train
- Did not generalise even within training domain
 - begin failing at n=39 (vs 257 for LSTM)
- Did not learn any discernible counting mechanism

Trained $a^n b^n c^n$, (on positive examples up to length 50)

Activations on $a^{100}b^{100}c^{100}$:



Trained $a^n b^n c^n$, (on positive examples up to length 100)

Activations on $a^{100}b^{100}c^{100}$:



GRU:

- Took much longer to train
- Did not generalise well
 - begin failing at n=9 (vs 101 for LSTM)
- Did not learn any discernible counting mechanism

Conclusion

IRNN LSTM SRNN GRU Trainability



Take Home Message

Architectural Choices Matter!

and result in actual differences in expressive power

Don't fall in the Turing Tarpit!

Thank You

GitHub repository:

https://github.com/tech-srl/counting_dimensions

Google Colab (link through GitHub as well):

https://tinyurl.com/ybjkumrz