EBMT Based on Finite Automata State Transfer Generation

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Introduction

- EBMT : a method of translation by the principle of analogy
- > Three basic modules
 - Matching module
 - Alignment module
 - Recombination module
- > The last two modules can be regarded as a translation generation module.
 - Semantic-based generation approach
 - > Obtains an appropriate translation fragment for each part of the input sentence.
 - > Final translation is generated by recombining the translation fragments in some order.
 - > Shortcoming: doesn't take into account the fluency between the translation fragments
 - Statistical approach
 - > Selects translation fragments with a statistical model
 - Can improve the fluency between the translation fragments by using n-gram co-occurrence statistics.
 - Shortcoming: doesn't take into account the semantic relation between the example and the input sentence
 - > Method based on tree string correspondence (TSC) and statistical generation
 - Can solve the shortcomings of the above generation approaches ;
 - But: depends on the tree parser so much that if the parser doesn't work well, it is impossible to generate a proper translation result.

System Structure of Our CJ EBMT System

- Our generation method
 - Uses the target sentence of the selected example to generate the translation of the input sentence.
 - Generate the translation in a finite automata state transfer manner.



Structure of our EBMT System

Generation Based on Finite Automata State Transfer

> Matching: select translation examples for the input sentence

Method: a combined method based on substantive word matching and stop word matching

WordSim(A,B)= $2 \cdot \frac{SameWC(A,B)}{len(A) + len(B)}$

StopWord_Sim(A,B) = exp(($abs(StopWord(A) - StopWord(B)) \times \beta$)

 $final _Sim(A,B) = \frac{WordSim(A,B)}{StopWord _Sim(A,B)}$

Generation

- Step 1 、 Build links from the fragments in the input sentence to the fragments in the target sentence of the selected example
- Step 2 , Assign states to each of these links
- Step 3 Construct a finite automaton and generate the translation result in an automaton state transfer manner

Step 1 for Generation: Building Links

- Link : a link from a fragment in one sentence S_1 to a fragment in another sentence S_2 is defined as a 3-tuple (Sf_i , Tf_j , t).
 - \succ Sf_i: a fragment in S₁
 - \succ Tf_i : a fragment in S_2
 - t: link type, we define four link types: I, R, D, N, which mean inserting, replacing, deleting and outputting directly respectively
- > Build links from the fragments in the input sentence S to the fragments in the target sentence B of the selected example (A, B)
 - First: Build links from S's fragments to A's fragments using a revised edit distance

algorithm (will be shown in the next slide). Its result is denoted as $LinkSet(S \rightarrow A)$.

- Second: Build links from S's fragments to B's fragments (denoted as $LinkSet(S \rightarrow B)$) according to following rules.
 - (a) For a link in *LinkSet(S→A)*, if neither its source fragment nor its target fragment is null, replace its target fragment with this target fragment's corresponding aligned fragment in *B*, and add this new link to *LinkSet(S→B)*.
 - (b) For a link in *LinkSet(S→A)* whose target fragment is null, add it to *LinkSet(S→B)* directly.
 - (c) For those fragments in B that have not been linked, build links for each of them by assigning a null source fragment and a D link type to them respectively, and add these links to LinkSet(S→B).
 - (d) Reorder the items of $LinkSet(S \rightarrow B)$ in their target fragments' order in sentence B Feiling Ren 2007-10

Step 1 for Generation: Building Links

The algorithm for building links from S's fragments to A's fragments is shown

```
as followings.
m = length(S_1), n = length(S_2)
d[0][0] = 0; tags[0][0] = 0;
for i=1 to m
   d[i][0]=q+d[i-1][0]; tags[i][0]='D'
for j=1 to n
   d[0][j]=r+d[0][j-1]; tags[0][j]='I'
for i=1 to m
  for j=1 to n
    p = computeCost(S_{i}[i-1], S_{2}[j-1]);
    a = d[i-1][j-1] + p;
    b=d[i-1][j] + q;
    c=d[i][j-1] + r;
    d[i][j] = min(a,b,c);
    if(min=a and p==0)
        tags[i][j] = N';
    else if (min=a)
        tags[i][j] = R';
     else if (min=b)
   Revised Edit Distance Algorithm
```

computeCost is a function to compute two fragments' linking cost based on their lexical forms and their head words' POSs.

- If two fragments' lexical forms are the same and their head words' POSs are the same too, this cost is zero;
- if two fragments' lexical forms are the same but their head words' POSs are different, this cost is 0.2;
- otherwise, this value is assigned by human's experiences according to the two fragments' head words' POSs as shown in the following table

$PosPair(c_i, c_j)$	W _i
(noun, noun)	0.5
(noun, auxiliary)	0.8
(noun, adjective)	0.85

Linking Cost for Two Fragments

Step 1 for Generation: Building Links

> The whole process of this step can be shown in the following figure



Step 2 for generation: States Assignment

- States for Non-I Type's Links
 - > If its link type is R, a state named S_R is assigned
 - \succ If its link type is *D*, a state named *S_D* is assigned;
 - \succ If its link type is *N*, a state named *S*_*N* is assigned.
- > States for *I* Type's Links
 - Consider context of current *I*-type link's pre- and post- links
 - Consider link shapes
 - Define 12 basic link shapes and 3 extended link shapes for *I*-type link, and map each of these link shapes to an *I*-type link's state.

Step 2 for generation: States Assignment

➢ Basic States for *I*-type's Link



Step 2 for generation: States Assignment



> Extended states can be converted into basic states

- For state 13, move rightward until find a non-*I* type's link, if this link's target fragment is null, convert it to state 6; otherwise, convert it to a state among state 1 to state 5 according to the link shapes of fragment *i*-1's link and the new found link; if can't find a non-*I* type's link in current link's right side, convert it to state 11.
- For state 14, move rightward until find a non-*I* type's link, if this link's target fragment is null, convert it to state 8, otherwise, convert it to state 7; if can't find a non-*I* type's link in current link's right side, convert it to state 12.
- For state 15, move rightward until find a non-*I* type's link, if this link's target fragment is null, convert it to state 10, otherwise, convert it to state 9; if can't find a non-*I* type's link in current link's right side, move leftward until find a non-*I* type's link (this link will be found always) and convert it to state 11.

Step 3 for generation: Translation Generation

- Generation Operation for Non-I Type Links' States
 - If a link's state is S_R, replace this link's target fragment with its source fragment's translation, and denote this operation as O(R);
 - If a link's state is S_D, delete this link's target fragment, and denotes this operation as O(D);
 - ➢ If a link's state is S_N, remain this link's target fragment unchanged, and denote this operation as O(N).
- Generation Operation for I Type Links' States
 - Take its source fragment's pre- and post- fragments into account and judge: whether the fragment combinations (*i*-1,*i*,*i*+1), (*i*-1,*i*) and (*i*,*i*+1) are chunks. If they are chunks, look up their corresponding translations in dictionary, otherwise, look up *i*'s translation in dictionary (we assume its translation can be found always).

According to current *I*-type link's state and the recognized chunk information, we choose one of these chunks as current *I*-type link's new source fragment for later processing, and define 10 possible generation

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Step 3 for generation: Translation Generation

➢ Generation Operation for *I* Type Links' States

- O(0): Delete the links that take *B*'s fragments among m+1 to *n* as their target fragments. And for the link that takes *B*'s fragment *m* as target fragment, replace *m* with the translation of current *I*-type link's new source fragment.
- **O(1)**: For the link that takes *B*'s fragment *m* as target fragment, replace *m* with the translation of current *I*-type link's new source fragment.
- *O(2):* For the link that takes *B*'s fragment *n* as target fragment, replace *n* with the translation of current *I*-type link's new source fragment.
- *O(3):* For the link that takes *B*'s fragment *m* as target fragment, add the translation of current *I*-type link's new source fragment to the end of *m*.
- *O(4):* For the link that takes *B*'s fragment *n* as target fragment, add the translation of current *I*-type link's new source fragment to the end of *n*.
- **O(5):** For the link that takes *B*'s fragment *m* as target fragment, replace *m* with the translation of current *I*-type link's new source fragment. And delete the link that takes *B*'s fragment *n* as target fragment.
- **O(6):** For the link that takes *B*'s fragment *n* as target fragment, replace *n* with the translation of current *I*-type link's new source fragment. And delete the link that takes *B*'s fragment *m* as target fragment.
- *O(7)*: For the link that takes *B*'s fragment *m* as target fragment, add the translation of current *I*-type link's new source fragment before *m*.
- **O(8):** For the link that takes *B*'s fragment *n* as target fragment, add the translation of current *I*-type link's new source fragment before *n*.
- **O(9):** Do not modify any link's target fragment.

Step 3 for generation: Translation Generation

▷ Based on $LinkSet(S \rightarrow B)$ and the assigned states, we construct an automaton that has a similar form as shown in the following figure



- \succ B is a start state
- \succ E is an end state
- \succ {*I*, *R*, *D*, *N*} are link types
- > $\{O(N), O(D), O(R)\}$ in parallelogram are the operations
- # is a fictitious symbol that indicates the end of the automaton's input
- {S_R, S_D, S_N} are states correspond to non-I type's links
- S_I' is a state set that corresponds to I-type's links

Step 3 for generation: Translation Generation



- O' in the operation of state 3 means the automaton generates the fragment combination (*i*-1,*i*,*i*+1)'s translation by simply joining their single fragment's translations together.
- d₁ means the semantic distance from fragment *i* to fragment
 i-1, and d₂ means the semantic distance from fragment *i* to fragment *i*+1, and they are computed as following formula:

$$dist(f_1, f_2) = \sum_{c_i \in f_1} \sum_{c_j \in f_2} w_k(PosPair(c_i, c_j))$$



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Step 3 for generation: An Example

- ▶ Suppose S is "他很爱他的妻子(He loves his wife very much)". The selected example (A,B) is "(他爱他的妈妈(He loves his mother),彼は、彼の母を愛し ています(He loves his mother))".
- ▶ After building links, *LinkSet(S → B)* is: (他(he), 彼(he), N), (null, は(ha), D), (很(very much), null, I), (他的(his), 彼の(his), N), (妻子(wife), 母(mother), R), (null, を(wo), D) (愛(loves), 愛しています(loves), N)
- ▶ Its corresponding state sequence is: S_N , S_D , S_I_4 (the forth state in the basic *I*-type's links), S_N , S_R , S_D , S_N .
- Construct an automaton, and begin to states transfer and t ranslation generation.
- For the link ($(\underline{h}(he), \underline{h}(he), N)$), its state is S_N . The automaton executes operation O(N) and does not mo dify this link's target fragment.
- For the link (null, (ha), D), its state is S_D. The automaton executes operation O(D) and deletes this link's target fragment.
- ▶ For the link (很(very much),null,I), its state is S_I_4. If the fragment combination (i-1,i) "他 很(he...very much)" is a chunk and the corresponding translation
- is "彼は、とても(he...very much)", the automaton executes operation O(1). It first takes this recognized chunk as current link's new source fragment. Then it
- selects the link whose target fragment is "彼(he)", and this link is (他(he),彼(he),N). Thirdly, it replaces the selected link's target fragment with the
- translation of current *I*-type link's new source fragment. At last the selected link is changed to (他(he), 彼は、とても(he…very much), N).
- ト For the link (他的(*his*),彼の母(*his*),N), its state is S_N . The automaton executes oper ation O(N) and does not modify this link's target fragment.
- ▶ For the link (妻子(*wife*),母(*mother*),R), its state is S_R. The automaton executes operation O(R) and replaces this link's target fragment with its source fragment's translation. Finally current link is changed to (妻子(*wife*),妻(*wife*),R).
- For the link (null, $\mathcal{E}(wo), D$), its state is S_D. The automaton executes operation O(D) and deletes this link's target fragment.
- ▶ For the link (爱(loves),愛しています(loves),N), its state is S_N. The automaton executes operation O(N) and does not modify this link's target fragment.
- ▶ At last, the automaton ends the state tran sfer process and outputs *LinkSet(S→B)*'s modified target fragment s equence "彼は、とても彼の妻愛しています(he loves his wife very much)" and takes it as the input sentence 's translation.

Experiments

System Resources

- Bilingual Corpus: We collect 10083 Chinese-Japanese bilingual sentences from Internet in Olympic domain as examples
- Bilingual Dictionary: A bilingual dictionary is used to translate the input fragment and to judge whether an input fragment is a chunk.
- Language Model: We collected an approximate 1,400,000 words' Japanese monolingual corpus and a similar size's Chinese monolingual corpus from Internet, and trained a standard trigram Japanese language model for Chinese-to-Japanese EBMT system and a standard trigram Chinese language model for Japanese-to-Chinese EBMT system respectively.
- Test Corpus: We collect another 100 bilingual sentences in Olympic domain from Internet as test corpus.

Experimental Result

Experimental Results for Chinese-to-Japanese EBMT

Suctam

Method	NIST	BLEU
Baseline	4.8321	0.4913
Our System	5.9729	0.7705

Experimental Results for Japanese-to-Chinese EBMT

S	vstem

	500111	
Method	NIST	BLEU
Baseline	4.1275	0.4076
Our System	5.0976	0.5908

Experiments----Some Translation Examples

Input:	我们的足球被对方前锋拦截
Output:	私 たち の サッカー は 相手 の 前鋒 に 阻まれ た
Input:	摔跤强国俄罗斯和日本有很多足球俱乐部
Output:	レスリング の 強国 ロシア と 日本 に は 很多 足球 俱乐部 が ある
Input:	中国运动员孙英杰今年一直主攻马拉松
Output:	中国 の スポーツ 選手 英傑 は 今年 ずっと マラソン を 専攻 として いる
	Some Translation Results for Chinese-to-Japanese Translation
Input:	審判員はいささかのためらいもなくペナルティーキックを科した
Output:	裁判 毫不犹豫 地 判罚 点球
Input:	中国のチームにはマンツーマンディフェンス戦術がある
Output:	中国 的 队 有 盯人 战术
Input:	スウェ - デンのチ - ム20分のペナルティを受けた
Output:	瑞典 队 被 罚 了 20 分

Some Translation Results for Japanese-to-Chinese Translation

Conclusions and Future Work

Conclusions:

- > The natural of the states are some transfer rules.
- > Our work can work on most of language pairs.
- It doesn't need any complicated parsers.
- Future Work
 - Merge syntax analysis into our method
 - > Merge probability knowledge into state assignment and generation.

The End

≻Thanks!

If you have any question, please contact me by renfeiliang@gmail.com, or renfeiliang@ise.neu.edu.cn

Welcome to my website: http://www.nlplab.cn/renfeiliang/

