

Constraint-Based Categorical Grammar

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Abstract

We propose a generalization of Categorical Grammar in which lexical categories are defined by means of recursive constraints. In particular, the introduction of relational constraints allows one to capture the effects of (recursive) lexical rules in a computationally attractive manner. We illustrate the linguistic merits of the new approach by showing how it accounts for the syntax of Dutch cross-serial dependencies and the position and scope of adjuncts in such constructions. Delayed evaluation is used to process grammars containing recursive constraints.

1 Introduction

Combinations of Categorical Grammar (CG) and unification naturally lead to the introduction of polymorphic categories. Thus, Karttunen (1989) categorizes NP's as X/X , where X is a verbal category, Zeevat *et al.* (1987) assign the category $X/(NP \setminus X)$ to NP's, and Emms (1993) extends the Lambek-calculus with polymorphic categories to account for coordination, quantifier scope, and extraction.

The role of polymorphism has been restricted, however, by the fact that in previous work categories were defined as feature structures using the simple, non-recursive, constraints familiar from feature description languages such as PATR. Relational constraints can be used to define a range of polymorphic categories that are beyond the expressive capabilities of previous approaches.

In particular, the introduction of relational constraints captures the effects of (recursive) lexical rules in a computationally attractive manner. The addition of such rules makes it feasible to consider truly 'lexicalist' grammars, in which a powerful lexical component is accompanied by a highly restricted syntactic component, consisting of application only.

2 Recursive Constraints

In CG, many grammatical concepts can only be defined recursively. Dowty (1982) defines grammatical functions such as *subject* and *object* as being the ultimate and penultimate 'argument-in' of a verbal category. Hoeksema (1984) defines verbs as exocentric categories reducible to s . Lexical rules frequently refer to such concepts. For instance, a categorial lexical rule of passive applies to *verbs* selecting an *object* and must remove the *subject*.

In standard unification-based formalisms, these concepts and the rules referring to such concepts cannot be expressed directly.

2.1 Subject-verb agreement

Consider a categorial treatment of subject-verb agreement with intransitive ($NP[NOM] \setminus s$) and transitive ($(NP[NOM] \setminus s)/NP[ACC]$) verbs defined as follows:

- (1) $lex(walks, X) :-$
 $iv(X).$
 $lex(kisses, X) :-$
 $tv(X).$

$$iv\left(\begin{array}{l} val [cat s] \\ dir '\setminus' \\ arg \left[\begin{array}{l} cat np \\ case nom \end{array} \right] \end{array}\right).$$

$$tv\left(\begin{array}{l} val \left[\begin{array}{l} val [cat s] \\ dir '\setminus' \\ arg \left[\begin{array}{l} cat np \\ case nom \end{array} \right] \end{array} \right] \\ dir '/' \\ arg \left[\begin{array}{l} cat np \\ case acc \end{array} \right] \end{array}\right).$$

Subject-verb agreement can be incorporated easily if one reduces agreement to a form of subcategorization.

If, however, one wishes to distinguish these two pieces of information (to avoid a proliferation of subcategorization types or for morphological reasons, for instance), it is not obvious how this could be done without recursive constraints. For intransitive verbs one needs the constraint that $\langle \text{arg agr} \rangle = \text{Agr}$ (where Agr is some agreement value), for transitive verbs that $\langle \text{val arg agr} \rangle = \text{Agr}$, and for ditransitive verbs that $\langle \text{val val arg agr} \rangle = \text{Agr}$. The generalization is captured using the recursive constraint *sv_agreement* (2). In (2) and below, we use definite clauses to define lexical entries and constraints. Note that lexical entries relate words to feature structures that are defined indirectly as a combination of simple constraints (evaluated by means of unification) and recursive constraints.¹

- (2) *lex*(walks, X) :-
 iv(X),
 sv_agreement(sg3, X).
lex(kisses, X) :-
 tv(X),
 sv_agreement(sg3, X).

sv_agreement(Agr, $\left[\begin{array}{l} \text{cat np} \\ \text{agr Agr} \end{array} \right] \backslash \text{S}$).
sv_agreement(Agr, Y \ X) :-
 sv_agreement(Agr, X).

Relational constraints can also be used to capture the effect of lexical rules. In a lexicalist theory such as CG, in which syntactic rules are considered to be universally valid scheme's of functor-argument combination, lexical rules are an essential tool for capturing language-specific generalizations. As Carpenter (1991) observes, some of the rules that have been proposed must be able to operate recursively. Predicative formation in English, for instance, uses a lexical rule turning a category reducible to VP into a category reducing to a VP-modifier (VP \ VP). As a VP-modifier is reducible to VP, the rule can (and sometimes must) be applied recursively.

2.2 Adjuncts as arguments

Miller (1992) proposes a lexical rule for French nouns which adds an (modifying) adjective to the list of arguments that the noun subcategorizes for. Since a noun

¹We use X/Y and $Y \ X$ as shorthand for $\left[\begin{array}{l} \text{val } X \\ \text{dir } '/' \\ \text{arg } Y \end{array} \right]$ and $\left[\begin{array}{l} \text{val } X \\ \text{dir } '\backslash' \\ \text{arg } Y \end{array} \right]$, respectively and S, NP, and Adj as 'typed variables' of type $\left[\text{cat } s \right]$, $\left[\text{cat } \text{np} \right]$, and $\left[\text{cat } \text{adj} \right]$, respectively.

can be modified by any number of adjectives, the rule must be optional as well as recursive. The advantages of using a lexical rule in this case is that it simplifies accounting for agreement between nouns and adjectives and that it enables an account of word order constraints between arguments and modifiers of a noun in terms of *obliqueness*.

The idea that modifiers are introduced by means of a lexical rule can be extended to verbs. That is, adjuncts could be introduced by means of a recursive rule that optionally adds these elements to verbal categories. Such a rule would be an alternative for the standard categorial analysis of adjuncts as (endocentric) functors. There is reason to consider this alternative.

In Dutch, for instance, the position of verb modifiers is not fixed. Adjuncts can in principle occur anywhere to the left of the verb:²

- (3) a. dat Johan opzettelijk een ongeluk
 that J. deliberately an accident
 veroorzaakt
 causes
 that J. deliberately causes an accident
 b. dat Johan Marie opzettelijk
 that J. M. deliberately
 geen cadeau geeft
 no present gives
 that J. deliberately gave M. no present

There are several ways to account for this fact. One can assign multiple categories to adjuncts or one can assign a polymorphic category x/x to adjuncts, with x restricted to 'verbal projections' (Bouma, 1988).

Alternatively, one can assume that adjuncts are not functors, but arguments of the verb. Since adjuncts are optional, can be iterated, and can occur in several positions, this implies that verbs must be polymorphic. The constraint *add_adjuncts* has this effect, as it optionally adds one or more adjuncts as arguments to the 'initial' category of a verb:

- (4) *lex*(veroorzaken, X) :-
 add_adjuncts(X, NP \ (NP \ S)).
lex(geven, X) :-
 add_adjuncts(X, NP \ (NP \ (NP \ S))).

add_adjuncts(S, S).
add_adjuncts(Adj \ X, Y) :-
 add_adjuncts(X, Y).
add_adjuncts($\left[\begin{array}{l} \text{val } X \\ \text{dir } D \\ \text{arg } A \end{array} \right]$, $\left[\begin{array}{l} \text{val } Y \\ \text{dir } D \\ \text{arg } A \end{array} \right]$) :-
 add_adjuncts(X, Y).

²As we want to abstract away from the effects of 'verb-second', we present only examples of subordinate clauses.

This constraint captures the effect of applying the following (schematic) lexical rule recursively:

$$(5) \quad X_1 \setminus \dots \setminus X_i \setminus X_{i+1} \setminus \dots \setminus S / Y_1 \dots Y_n \\ \Downarrow \\ X_1 \setminus \dots \setminus X_i \setminus \text{Adj} \setminus X_{i+1} \setminus \dots \setminus S / Y_1 \dots Y_n$$

The derivation of (3a) is given below (where $X \Rightarrow Y$ indicates that $\text{add_adjuncts}(Y, X)$ is satisfied, and $\text{IV} = \text{NP} \setminus \text{S}$).

$$(6) \quad \begin{array}{cccc} \dots J. & \text{opzettelijk} & \text{een ongeluk} & \text{veroorzaakt} \\ \text{NP} & \text{ADJ} & \text{NP} & \text{NP} \setminus \text{IV} \\ & & & \Downarrow \\ & & & \text{NP} \setminus (\text{ADJ} \setminus \text{IV}) \\ & & \text{ADJ} \setminus \text{IV} & \\ \hline & \text{IV} & & \\ \hline \text{S} & & & \end{array}$$

An interesting implication of this analysis is that in a categorial setting the notion ‘head’ can be equated with the notion ‘main functor’. This has been proposed by Barry and Pickering (1990), but they are forced to assign a category containing Kleene-star operators to verbal elements. The semantic counterpart of such category-assignments is unclear. The present proposal is an alternative for such assignments which avoids introducing new categorial operators and which does not lead to semantic complications (the semantics of add_adjuncts is presented in section 3.3). Below we argue that this analysis also allows for a straightforward explanation of the distribution and scope of adjuncts in verb phrases headed by a verbal complex.

3 Cross-Serial Dependencies

In Dutch, verbs selecting an infinitival complement (e.g. modals and perception verbs) give rise to so called cross-serial dependencies. The arguments of the verbs involved appear in the same order as the verbs in the ‘verb cluster’:

- (7) a. dat An₁ Bea₂ wil₁ kussen₂.
 dat An Bea wants to kiss
that An wants to kiss Bea
 b. dat An₁ Bea₂ Cor₃ wil₁
 dat An Bea Cor wants
 zien₂ kussen₃.
 to see kiss
that An wants to see Bea kiss Cor

The property of forming cross-serial dependencies is a lexical property of the matrix verb. If this verb is a ‘trigger’ for cross-serial word order, this order is obligatory, whereas if it is not, the infinitival complement will

follow the verb:

- (8) a. *dat An wil Bea kussen.
 b. dat An zich voornam Bea
 that An Refl. planned Bea
 te kussen.
 to kiss
that An. planned to kiss Bea
 c. *dat An zich Bea voornam te kussen.

3.1 Generalized Division

Categorial accounts of cross-serial dependencies initially made use of a syntactic rule of composition (Steedman, 1985). Recognizing the lexical nature of the process, more recent proposals have used either a lexical rule of composition (Moortgat, 1988) or a lexical rule of ‘division’ (Hoeksema, 1991). Division is a rule which enables a functor to inherit the arguments of its argument:³

$$X/Y \Rightarrow (X/Z_1 \dots /Z_n)/(Y/Z \dots /Z_n)$$

To generate cross-serial dependencies, a ‘disharmonic’ version of this rule is needed:

$$(9) X/Y \Rightarrow (Z_1 \setminus \dots Z_n \setminus X)/(Z_1 \setminus \dots Z_n \setminus Y)$$

Hoeksema proposes that verbs which trigger cross-serial word order are subject to (9):

$$(10) \quad \begin{array}{cccc} \dots An & Bea & wil & kussen \\ \text{NP} & \text{NP} & \text{IV} / \text{IV} & \text{NP} \setminus \text{IV} \\ & & \Downarrow & \\ & & (\text{NP} \setminus \text{IV}) / (\text{NP} \setminus \text{IV}) & \\ \hline & & \text{NP} \setminus \text{IV} & \\ \hline \text{IV} & & & \end{array}$$

In a framework using recursive constraints, generalized disharmonic division can be implemented as a recursive constraint connecting the initial category of such verbs with a derived category:

- (11) $\text{lex}(\text{willen}, X) :-$
 $\text{cross_serial}(X, (\text{NP} \setminus \text{S}) / (\text{NP} \setminus \text{S})).$
 $\text{lex}(\text{zien}, X) :-$
 $\text{cross_serial}(X, (\text{NP} \setminus (\text{NP} \setminus \text{S})) / (\text{NP} \setminus \text{S})).$
 $\text{lex}(\text{voornemen}, (\text{NP}_{\text{refl}} \setminus (\text{NP} \setminus \text{S})) / (\text{NP} \setminus \text{S})).$

³Argument inheritance is used in HPSG to account for verb clustering in German (Hinrichs and Nakazawa, 1989). The HPSG analysis is essentially equivalent to Hoeksema’s account.

(12) *cross_serial*(Out, In) :-
 division(Out, In),
 verb_cluster(Out).

division(X, X).
division((Z\X)/(Z\Y), X'/Y') :-
 division(X/Y, X'/Y').

verb_cluster([arg [vc +]]).

Only verbs that trigger the cross-serial order are subject to the *division* constraint. This accounts immediately for the fact that cross-serial orders do not arise with all verbs selecting infinitival complements.

3.2 Verb Clusters

The *verb_cluster* constraint ensures that cross-serial word order is obligatory for verbs subject to *cross_serial*. To rule out the ungrammatical (8a), for instance, we assume that *Bea kussen* is not a verb cluster. The verb *kussen* by itself, however, is unspecified for VC, and thus (7a) is not excluded.

We do not assume that cross-serial verbs take lexical arguments (as has sometimes been suggested), as that would rule out the possibility of complex constituents to the right of cross-serial verbs altogether. If one assumes that a possible bracketing of the verb cluster in (7b) is [*wil [zien kussen]*] (coordination and fronting data have been used as arguments that this is indeed the case), a cross-serial verb must be able to combine with non-lexical verb clusters. Furthermore, if a verb selects a particle, the particle can optionally be included in the verb cluster, and thus can appear either to the right or to the left of a governing cross-serial verb. For a verb cluster containing two cross-serial verbs, for instance, we have the following possibilities:

- (13) a. dat An Bea heeft durven aan
 that An Bea has dared part.
 te spreken
 to speak
 that An has dared to speak to Bea.
 b. dat An Bea heeft aan durven te spreken.
 c. dat An Bea aan heeft durven te spreken.

A final piece of evidence for the fact that cross-serial verbs may take complex phrases as argument stems from the observation that certain adjectival and prepositional arguments can also appear as part of the verb cluster:

- (14) dat An dit aan Bea had duidelijk
 that An this to Bea has clear
 gemaakt
 made
 that An had made this clear to Bea

Cross-serial verbs select a +VC argument. Therefore, all phrases that are not verb clusters must be marked -VC. In general, in combining a (verbal) functor with its argument, it is the argument that determines whether the resulting phrase is -VC. For instance, NP-arguments always give rise to -VC phrases, whereas particles and verbal arguments do not give rise to -VC phrases. This suggests that NP's must be marked -VC, that particles and verbs can remain unspecified for this feature, and that in the syntactic rule for application the value of the feature VC must be reentrant between argument and resultant.

3.3 The distribution and scope of adjuncts

The analysis of cross-serial dependencies in terms of argument inheritance interacts with the analysis of adjuncts presented in section 2.2. If a matrix verb inherits the arguments of the verb it governs, it should be possible to find modifiers of the matrix verb between this verb and one of its inherited arguments. This prediction is borne out (15a). However, we also find structurally similar examples in which the adjunct modifies the governed verb (15b). Finally, there are examples that are ambiguous between a wide and narrow scope reading (15c). We take it that the latter case is actually what needs to be accounted for, i.e. examples such as (15a) and (15b) are cases in which there is a strong preference for a wide and narrow scope reading, respectively, but we will remain silent about the (semantic) factors determining such preferences.

- (15) a. dat Frits Marie volgens mij lijkt
 that F. M. to me seems
 te ontwijken.
 to avoid
 It seems to me that F. avoids M.
 b. dat Frits Marie opzettelijk lijkt
 that F. M. deliberately seems
 te ontwijken.
 to avoid
 It seems that F. deliberately avoids M.
 c. dat Frits Marie de laatste tijd lijkt
 that F. M. lately seems
 te ontwijken.
 to avoid
 It seems lately as if F. avoids M.
 It seems as if F. avoids M. lately

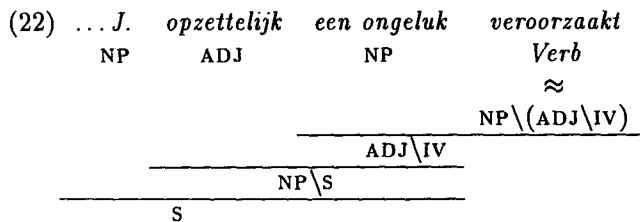
On the assumption that the lexical entries for *lijken* en *ontwijken* are as in (16), example (15c) has two possible derivations ((17) and (18)). Procedurally speaking, the rule that adds adjuncts can be applied either to the matrix verb (after division has taken place) or to the

considers lexical rules as unary syntactic rules, to non-branching derivations of unbounded length. In both cases, a parser may not terminate. One of the main advantages of modeling lexical rules by means of constraints is that it suggests a solution for this problem. A control strategy which delays the evaluation of constraints until certain crucial bits of information are filled in avoids non-termination and in practice leads to grammars in which all constraints are fully evaluated at the end of the parse-process.

Consider a grammar in which the only recursive constraint is *add_adjuncts*, as defined in section 2.2. The introduction of recursive constraints in itself does not solve the non-termination problem. If all solutions for *add_adjuncts* are simply enumerated during lexical look-up an infinite number of categories for any given verb will result.

During processing, however, it is not necessarily the case that we need to consider all solutions. Syntactic processing can lead to a (partial) instantiation of the arguments of a constraint. If the right pieces of information are instantiated, the constraint will only have a finite number of solutions.

Consider, for instance, a parse for the following string.



Even if the category of the verb is left completely open initially, there is only one derivation for this string that reduces to S (remember that the syntax uses application only). This derivation provides the information that the variable *Verb* must be a transitive verb selecting one additional adjunct, and with this information it is easy to check whether the following constraint is satisfied:

add_adjuncts(NP \ (ADJ \ (NP \ S)), NP \ (NP \ S)).

This suggests that recursive constraints should not be evaluated during lexical look-up, but that their evaluation should be delayed until the arguments are sufficiently instantiated.

To implement this delayed evaluation strategy, we used the **block** facility of Sicstus Prolog. For each recursive constraint, a **block** declaration defines what the conditions are under which it may be evaluated. The definition of *add_adjuncts* (with semantics omitted for

readability), for instance, now becomes:

(23) *add_adjuncts*([*arg Arg*]_X, Y) :-
add_adjuncts(X, Y, Arg).

:- **block** *add_adjuncts*(?, ?, -).

add_adjuncts(S, S, -).

add_adjuncts(Adj \ X, Y, -) :-

add_adjuncts(X, Y).

add_adjuncts($\begin{bmatrix} \text{val } X \\ \text{dir } D \\ \text{arg } A \end{bmatrix}$, $\begin{bmatrix} \text{val } Y \\ \text{dir } D \\ \text{arg } A \end{bmatrix}$, A) :-

add_adjuncts(X, Y).

We use *add_adjuncts/2* to extract the information that determines when *add_adjuncts/3* is to be evaluated. The **block** declaration states that *add_adjuncts/3* may only be evaluated if the third argument (i.e. the argument of the 'output' category) is not a variable. During lexical look-up, this argument is uninstantiated, and thus, no evaluation takes place. As soon as a verb combines with an argument, the argument category of the verb is instantiated and *add_adjuncts/3* will be evaluated. Note, however, that calls to *add_adjuncts/3* are recursive, and thus one evaluation step may lead to another call to *add_adjuncts/3*, which in its turn will be blocked until the argument has been instantiated sufficiently. Thus, the recursive constraint is evaluated incrementally, with each syntactic application step leading to a new evaluation step of the blocked constraint. The recursion will stop if an atomic category s is found.

Delayed evaluation leads to a processing model in which the evaluation of lexical constraints and the construction of derivational structure is completely intertwined.

4.1 Other strategies

The delayed evaluation techniques discussed above can be easily implemented in parsers which rely on backtracking for their search. For the grammars that we have worked with, a simple bottom-up (shift-reduce) parser combined with delayed evaluation guarantees termination of the parsing process.

To obtain an efficient parser more complicated search strategies are required. However, chart-based search techniques are not easily generalized for grammars which make use of complex constraints. Even if the theoretical problems can be solved (Johnson, 1993; Dörre, 1993) severe practical problems might surface, if the constraints are as complex as the ones proposed here.

As an alternative we have implemented chart-based parsers using the 'non-interleaved pruning' strategy (terminology from (Maxwell III and Kaplan, 1994)).

Using this strategy the parser first builds a parse-forest for a sentence on the basis of the context-free backbone of the grammar. In a second processing phase parses are recovered on the basis of the parse forest and the corresponding constraints are applied. This may be advantageous if the context-free backbone of the grammar is 'informative' enough to filter many unsuccessful partial derivations that the parser otherwise would have to check.

As clearly a CUG grammar does not contain such an informative context-free backbone a further step is to use 'selective feature movement' (cf. again (Maxwell III and Kaplan, 1994)). In this approach the base grammar is compiled into an equivalent modified grammar in which certain constraints from the base grammar are converted to a more complex context-free backbone in the modified grammar.

Again, this technique does not easily give good results for grammars of the type described. It is not clear at all where we should begin extracting appropriate features for such a modified grammar, because most information passing is simply too 'indirect' to be easily compiled into a context-free backbone.

We achieved the best results by using a 'hand-fabricated' context-free grammar as the first phase of parsing. This context-free grammar builds a parse forest that is then used by the 'real' grammar to obtain appropriate representation(s) for the input sentence. This turned out to reduce parsing times considerably.

Clearly such a strategy raises questions on the relation between this context-free grammar and the CUG grammar. The context-free grammar is required to produce a superset of the derivations allowed by the CUG. Given the problems mentioned above it is difficult to show that this is indeed the case (if it were easy, then it probably would also be easy to obtain such a context-free grammar automatically).

The strategy can be described in somewhat more detail as follows. The context-free phase of processing builds a number of items defining the parse forest, in a format that can be used by the second processing phase. Such items are four-tuples

$$\langle R, P_0, P, D \rangle$$

where R is a rule name (consistent with the rule names from the CUG), P_0, P are string positions and D describes the string positions associated with each daughter of the rule (indicating which part of the string is covered by that daughter).

Through a head-driven recursive descent the second processing phase recovers derivations on the basis of these items. Note that the delayed evaluation technique for complex constraints is essential here. Alternative solutions are obtained by backtracking. If the

first phase has done a good job in pruning many failing search branches then this is not too expensive, and we do not have to worry about the interaction of caching and complex constraints.

5 Final Remarks

In sections 2 and 3 we have sketched an analysis of cross-serial dependency constructions and its interaction with the position and scope of adjuncts. The rules given there are actually part of a larger fragment that covers the syntax of Dutch verb clusters in more detail. The fragment accounts for cross-serial dependencies and extraposition constructions (including cases of 'partial' extraposition), *infinitivus-pro-participio*, modal and participle inversion, the position of particles in verb clusters, clitic climbing, partial VP-topicalization, and verb second. In the larger fragment, additional recursive constraints are introduced, but the syntax is still restricted to application only.

The result of Carpenter (1991) emphasizes the importance of lexical rules. There is a tendency in both CG and HPSG to rely more and more on mechanisms (such as inheritance and lexical rules or recursive constraints) that operate in the lexicon. The unrestricted generative capacity of recursive lexical rules implies that the remaining role of syntax can be extremely simple. In the examples above we have stressed this by giving an account for the syntax of cross-serial dependencies (a construction that is, given some additional assumptions, not context-free) using application only. In general, such an approach seems promising, as it locates the sources of complexity for a given grammar in one place, namely the lexicon.

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