Rational Recurrences

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Abstract

Despite the tremendous empirical success of neural models in natural language processing, many of them lack the strong intuitions that accompany classical machine learning approaches. Recently, connections have been shown between convolutional neural networks (CNNs) and weighted finite state automata (WFSAs), leading to new interpretations and insights. In this work, we show that some recurrent neural networks also share this connection to WFSAs. We characterize this connection formally, defining rational recurrences to be recurrent hidden state update functions that can be written as the Forward calculation of a finite set of WFSAs. We show that several recent neural models use rational recurrences. Our analysis provides a fresh view of these models and facilitates devising new neural architectures that draw inspiration from WFSAs. We present one such model, which performs better than two recent baselines on language modeling and text classification. Our results demonstrate that transferring intuitions from classical models like WFSAs can be an effective approach to designing and understanding neural models.

1 Introduction

Neural models, and in particular gated variants of recurrent neural networks (RNNs, e.g., Hochreiter and Schmidhuber, 1997; Cho et al., 2014), have become a core building block for stateof-the-art approaches in NLP (Goldberg, 2016). While these models empirically outperform classical NLP methods on many tasks (Zaremba et al., 2014; Bahdanau et al., 2015; Dyer et al., 2016; Peng et al., 2017, *inter alia*), they typically lack the intuition offered by classical models, making it hard to understand the roles played by each of their components. In this work we show that many neural models are more interpretable than previously



Figure 1: A two-state WFSA \mathcal{B} described in §2. It is closely related to several models studied in this paper (§4.1). Bold circles indicate initial states, and double circles final states, which are associated with final weights. Arrows represent transitions, labeled by the symbols α they consume, and the weights as a function of α . Arcs not drawn are assumed to have weight $\overline{0}$. For brevity, $\forall \alpha$ means $\forall \alpha \in \Sigma$, with Σ being the alphabet.

thought, by drawing connections to weighted finite state automata (WFSAs). We study several recently proposed RNN architectures and show that one can use WFSAs to characterize their recurrent updates. We call such models **rational recurrences** (§3).¹ Analyzing recurrences in terms of WFSAs provides a new view of existing models and facilitates the development of new ones.

In recent work, Schwartz et al. (2018) introduced SoPa, an RNN constructed from WFSAs, and thus rational by our definition. They also showed that a single-layer max-pooled CNN (Le-Cun, 1998) can be simulated by a set of simple WFSAs (one per output dimension), and accordingly are also rational. In this paper we broaden such efforts, and show that rational recurrences are in frequent use (Mikolov et al., 2014; Balduzzi and Ghifary, 2016; Lei et al., 2016, 2017a,b; Bradbury et al., 2017; Foerster et al., 2017). For instance, we will show in §4 that the WFSA diagrammed

¹ Where the term *regular* is used with unweighted FSAs (e.g., regular languages, regular expressions), *rational* is the weighted analog (e.g., rational series, Sakarovitch, 2009; rational kernels, Cortes et al., 2004).

in Figure 1 has strong connections to several of the models mentioned above.

Based on these observations, we then discuss potential approaches to deriving novel neural architectures from WFSAs ($\S5$). As a case study, we present a new model motivated by the interpolation of a two-state WFSA and a threestate one, capturing (soft) unigram and bigram features, respectively. Our experiments show that in two tasks-language modeling and text classification-the proposed model outperforms recently proposed rational models ($\S6$). Further extensions might lead to larger gains, and the rational recurrence view could facilitate easier exploration of such extensions. To promote such exploration, we publicly release our implementation at https://github.com/ Noahs-ARK/rational-recurrences.

2 Background: Weighted Finite State Automata (WFSAs)

This section reviews weighted finite-state automata and semirings, which underly our analyses in §3. WFSAs extend nondeterministic *unweighted* finite-state automata by assigning weights to transitions, start states, and final states. Instead of simply accepting or rejecting a string, a WFSA returns a score for the string, and this score summarizes the weights along all paths through the WFSA that consume the string. In order for this summary score to be efficiently computable, weights are taken from a **semiring**.

Definition 1 (Kuich and Salomaa, 1986). A **semiring** is a set \mathbb{K} along with two associative binary operations on \mathbb{K} , \oplus (addition) and \otimes (multiplication), and two identity elements: $\overline{0}$ for addition, and $\overline{1}$ for multiplication. Semirings also require that addition is commutative, multiplication distributes over addition, and that multiplication by $\overline{0}$ annihilates (i.e., $\overline{0} \otimes a = a \otimes \overline{0} = \overline{0}$).

One common semiring is the **real** (or plustimes) semiring: $\langle \mathbb{R}, +, \cdot, 0, 1 \rangle$. The other one used in this work is the **max-plus** semiring $\langle \mathbb{R} \cup \{-\infty\}, \max, +, -\infty, 0 \rangle$. We refer the reader to Kuich and Salomaa (1986) for others.

Definition 2. A weighted finite-state automaton (WFSA) over a semiring \mathbb{K} is a 5-tuple, $\mathcal{A} = \langle \Sigma, \mathcal{Q}, \tau, \lambda, \rho \rangle$,² with:

- a finite input alphabet Σ ;
- a finite state set Q;
- transition weights $\tau : \mathcal{Q} \times \mathcal{Q} \times (\Sigma \cup \{\varepsilon\}) \to \mathbb{K};$
- initial weights $\lambda : \mathcal{Q} \to \mathbb{K}$;
- and final weights $\rho : \mathcal{Q} \to \mathbb{K}$.

 $\varepsilon \notin \Sigma$ marks special ε -transitions that may be taken without consuming any input. \mathcal{A} assigns a score $\mathcal{A}[\![\mathbf{x}]\!]$ to a string $\mathbf{x} = x_1 \dots x_n \in \Sigma^*$ by summing over the scores of all possible paths deriving \mathbf{x} . The score of each individual path is the product of the weights of the transitions it consists of. Formally:

Definition 3 (path score). Let $\pi = \pi_1 \dots \pi_n$ be a sequence of adjacent **transitions** in \mathcal{A} , with each transition $\pi_i = (q_i, q_{i+1}, z_i) \in \mathcal{Q} \times \mathcal{Q} \times (\Sigma \cup \{\varepsilon\})$. The path π **derives** string $\mathbf{x} \in \Sigma^*$, which is the substring of $\mathbf{z} = z_1 z_2 \dots z_n$ that excludes ε symbols (for example, if $\mathbf{z} = a\varepsilon bc\varepsilon\varepsilon\varepsilon d$, then $\mathbf{x} = abcd$). π 's score in \mathcal{A} is given by

$$\mathcal{A}[\boldsymbol{\pi}] = \lambda(q_1) \otimes \left(\bigotimes_{i=1}^n \tau(\pi_i)\right) \otimes \rho(q_{n+1}). \quad (1)$$

Definition 4 (string score). Let $\Pi(\mathbf{x})$ denote the set of all paths in \mathcal{A} that derive \mathbf{x} . Then the score assigned by \mathcal{A} to \mathbf{x} is defined to be

$$\mathcal{A}[\![\mathbf{x}]\!] = \bigoplus_{\boldsymbol{\pi} \in \Pi(\mathbf{x})} \mathcal{A}[\boldsymbol{\pi}].$$
(2)

Because \mathbb{K} is a semiring, $\mathcal{A}[\![\mathbf{x}]\!]$ can be computed in time linear in $|\mathbf{x}|$ by the Forward algorithm (Baum and Petrie, 1966). Here, for simplicity, we describe the Forward algorithm without ε -transitions.³ Its dynamic program is given by:

$$\Omega_0(q) = \lambda(q), \tag{3a}$$

$$\Omega_{i+1}(q) = \bigoplus_{q' \in \mathcal{Q}} \Omega_i(q') \otimes \tau(q', q, x_i), \quad (3b)$$

$$\mathcal{A}\llbracket \mathbf{x} \rrbracket = \bigoplus_{q \in \mathcal{Q}} \Omega_n(q) \otimes \rho(q).$$
(3c)

 $\Omega_i(q)$ gives the total score of all paths that derive $x_1 \dots x_i$ and end in state q.

Example 5. Figure 1 diagrams a WFSA \mathcal{B} , consisting of two states. A path starts from the initial state q_0 (with $\lambda(q_0) = \overline{1}$); it then takes any num-

²Some authors define τ , λ , and ρ to be partial functions applying only a subset of transitions, initial states, and final

states respectively. Our definition is equivalent, giving the weight functions value $\bar{0}$ wherever they were undefined.

 $^{{}^{3}\}varepsilon$ -transitions can be handled with a slight modification (Schwartz et al., 2018). Note though that if \mathcal{A} contains a cycle of ε -transitions, then either \mathbb{K} must follow the **star semiring** laws (Kuich and Salomaa, 1986), or the number of consecutive ε -transitions allowed must be capped.

ber of "self-loop" transitions, each consuming an input without changing the path score (since it's weighted by $\overline{1}$); it then consumes an input symbol α and takes a transition weighted by $\mu(\alpha)$, and reaches the final state q_1 (with $\rho(q_1) = \overline{1}$); it may further consume more input by taking selfloops at q_1 , updating the path score by multiplying it by $\phi(\alpha)$ for each symbol α . Then from Definition 4, we can calculate that \mathcal{B} gives the empty string score $\overline{0}$, and gives any nonempty string $\mathbf{x} = x_1 \dots x_n \in \Sigma^+$ score $\mathcal{B}[\![\mathbf{x}]\!] =$

$$\bigoplus_{i=1}^{n-1} \left(\mu(x_i) \otimes \bigotimes_{j=i+1}^n \phi(x_j) \right) \oplus \mu(x_n).$$
 (4)

B can be seen as capturing soft unigram patterns (Davidov et al., 2010), in the sense that it consumes one input symbol to reach the final state from the initial state. It is straightforward to design WFSAs capturing longer patterns by including more states (Schwartz et al., 2018), as we will discuss later in $\S4$ and $\S5$.

3 Rational Recurrences

Before formally defining rational recurrences in $\S3.2$, we highlight the connection between WFSAs and RNNs using a motivating example ($\S3.1$).

3.1 A Motivating Example

We describe a simplified RNN which strips away details of some recent RNNs, in order to highlight the behaviors of the forget gate and the input.

Example 6. For an input sequence $\mathbf{x} = x_1 \dots x_n$, let the word embedding vector for x_t be \mathbf{v}_t . As in many gated RNN variants (Hochreiter and Schmidhuber, 1997; Cho et al., 2014), we use a forget gate \mathbf{f}_t , which is computed with an affine transformation followed by an elementwise sigmoid function $\boldsymbol{\sigma}$. The current input representation \mathbf{u}_t is similarly computed, but with an optional nonlinearity (e.g., tanh) \boldsymbol{g} . The hidden state \mathbf{c}_t can be seen as a weighted sum of the previous state and the new input, controlled by the forget gate.

$$\mathbf{f}_t = \boldsymbol{\sigma} \big(\mathbf{W}_f \mathbf{v}_t + \mathbf{b}_f \big), \tag{5a}$$

$$\mathbf{u}_t = (\mathbf{1} - \mathbf{f}_t) \odot \boldsymbol{g} \big(\mathbf{W}_u \mathbf{v}_t + \mathbf{b}_u \big), \qquad (5b)$$

$$\mathbf{c}_t = \mathbf{f}_t \odot \mathbf{c}_{t-1} + \mathbf{u}_t. \tag{5c}$$

The hidden state \mathbf{c}_t can then be used in downstream computation, e.g., to calculate output state $\mathbf{h}_t = \tanh(\mathbf{c}_t)$, which is then fed to an MLP classifier. We focus only on the recurrent computation. In Example 6, both \mathbf{f}_t and \mathbf{u}_t depend only on the current input token x_t (through \mathbf{v}_t), and not the previous state. Importantly, the interaction with the previous state \mathbf{c}_{t-1} is *not* via affine transformations followed by nonlinearities, as in, e.g., an Elman network (Elman, 1990), where $\mathbf{c}_t =$ $tanh(\mathbf{W}_c \mathbf{c}_{t-1} + \mathbf{W}_v \mathbf{v}_t + \mathbf{b}_c)$. As we will discuss later, this is important in relating this recurrent update function to WFSAs.

Since the recurrent update in Equation 5c is elementwise, for simplicity we focus on just the *i*th dimension. Unrolling it in time steps, we get

$$[\mathbf{c}_t]_i = [\mathbf{f}_t]_i [\mathbf{c}_{t-1}]_i + [\mathbf{u}_t]_i$$
$$= \sum_{j=1}^{t-1} \left([\mathbf{u}_j]_i \prod_{k=j+1}^t [\mathbf{f}_k]_i \right) + [\mathbf{u}_t]_i,$$
(6)

where $[\cdot]_i$ denotes the *i*th dimension of a vector. As noted by Lee et al. (2017), the hidden state at time step *t* can be seen as a sum of previous input representations, weighted by the forget gate; longer histories typically get a smaller weight, since the forget gate values are between 0 and 1 due to the sigmoid function.

Let's recall the WFSA \mathcal{B} (Figure 1 and Example 5) using the real semiring $\langle \mathbb{R}, +, \cdot, 0, 1 \rangle$. Equation 6 is recovered by parameterizing \mathcal{B} 's weight functions μ and ϕ with

$$\mu(x_t) = [\mathbf{u}_t]_i, \quad \phi(x_t) = [\mathbf{f}_t]_i. \tag{7}$$

Denote the resulting WFSA by \mathcal{B}_i , and we have:

Proposition 7. Running a single layer RNN in Example 6 over any nonempty input string $\mathbf{x} \in \Sigma^+$, the *i*th dimension of *i*ts hidden state at time step t equals the score assigned by \mathcal{B}_i to $\mathbf{x}_{:t}$:

$$[\mathbf{c}_t]_i = \mathcal{B}_i[\![\mathbf{x}_{:t}]\!]. \tag{8}$$

In other words, the *i*th dimension of the RNN in Example 6 can be seen as a WFSA structurally equivalent to \mathcal{B} . Its weight functions are implemented as the *i*th dimension of Equations 5, and the learned parameters are the *i*th row of **W** and b. Then it is straightforward to recover the full *d*dimensional RNN, by collecting *d* such WFSAs, each of which is parametrized by a row in the **W**s and bs. Based on this observation, we are now ready to formally define rational recurrences.

3.2 Recurrences and Rationality

For a function $\mathbf{c}: \Sigma^* \to \mathbb{K}^d$, its *recurrence* is said to be the dependence of $\mathbf{c}(\mathbf{x}_{:t})$ on $\mathbf{c}(\mathbf{x}_{:t-1})$, for input

sequence $\forall \mathbf{x} \in \Sigma^+$. We discuss a class of recurrences that can be characterized by WFSAs. The mathematical counterpart of WFSAs are *rational power series* (Berstel and Reutenauer, 1988), justifying naming such recurrences *rational*:

Definition 8 (rational recurrence). The recurrence of $\mathbf{c} \colon \Sigma^* \to \mathbb{K}^d$ is said to be *rational*, if there exists a set of weighted finite state automata $\{\mathcal{A}_i\}_{i=1}^d$ over alphabet Σ and semiring $\langle \mathbb{K}, \oplus, \otimes, \overline{0}, \overline{1} \rangle$ with both \oplus and \otimes taking constant time and space, such that $\forall \mathbf{x} \in \Sigma^*$,

$$\left[\mathbf{c}(\mathbf{x})\right]_{i} = \mathcal{A}_{i}[\![\mathbf{x}]\!], \ \forall i \in \{1, 2, \dots, d\}.^{4}$$
(9)

It directly follows from Proposition 7 that

Corollary 9. *The recurrence in Example 6 is rational.*

4 Relationship to Existing Neural Models

This section studies several recently proposed neural architectures, and relates them to rational recurrences. §4.1 begins by relating some of them to the RNN defined in Example 6, and then to the WFSA \mathcal{B} (Example 5). We then describe a WFSA similar to \mathcal{B} , but with one additional state, and discuss how it provides a new view of RNN models motivated by *n*-gram features (§4.2). In §4.3 we study rational recurrences that are not elementwise, using an existing model.

In the following discussion, we shall assume the real seimiring, unless otherwise noted.

4.1 Neural Architectures Related to \mathcal{B}

Despite its simplicity, Example 6 corresponds to several existing neural architectures. For instance, quasi-RNN (QRNN; Bradbury et al., 2017) and simple recurrent unit (SRU; Lei et al., 2017b) aim to speed up the recurrent computation. To do so, they drop the matrix multiplication dependence on the previous hidden state, resulting in similar recurrences to that in Example 6.⁵ Other works start from different motivations, but land on similar recurrences, e.g., strongly-typed RNNs (T-RNN; Balduzzi and Ghifary, 2016) and its gated variants (T-LSTM and T-GRU), and structurally constrained RNNs (SCRN; Mikolov et al., 2014).

The analysis in §3.1 directly applies to SRU, T-RNN, and SCRN. In fact, Example 6 presents a slightly more complicated version of them. In these models, input representations are computed without the bias term or any nonlinearity: $\mathbf{u}_t = \mathbf{W}_u \mathbf{v}_t$. By Proposition 7 and Corollary 9:

Corollary 10. *The recurrences of single-layer SRU, T-RNN, and SCRN architectures are rational.*

It is slightly more complicated to analyze the recurrences of the QRNN, T-LSTM, and T-GRU. Although their hidden states c_t are updated in the same way as Equation 5c, the input representations and gates may depend on *previous inputs*. For example, in T-LSTM and T-GRU, the forget gate is a function of two consecutive inputs:

$$\mathbf{f}_t = \boldsymbol{\sigma} \left(\mathbf{V}_f \mathbf{v}_{t-1} + \mathbf{W}_f \mathbf{v}_t + \mathbf{b}_f \right). \tag{10}$$

QRNNs are similar, but may depend on up to K tokens, due to the K-window convolutions. Eisner (2002) discuss finite state machines for second (or higher) order probabilistic sequence models. Following the same intuition, we sketch the construction of WFSAs corresponding to QRNNs with 2-window convolutions in Appendix A, and summarize the key results here:

Proposition 11. The recurrences of single-layer *T*-GRU, *T*-LSTM, and QRNN are rational. In particular, a single-layer d-dimensional QRNN using K-window convolutions can be recovered by a set of d WFSAs, each with $O(2 |\Sigma|^{K-1})$ states.

The size of WFSAs needed to recover QRNN grows exponentially in the window size. Therefore, at least for QRNNs, Proposition 11 has more conceptual value than practical.

4.2 More than Two States

So far our discussion has centered on \mathcal{B} , a twostate WFSA capturing unigram patterns (Example 5). In the same spirit as going from unigram to *n*-gram features, one can use WFSAs with more states to capture longer patterns (Schwartz et al., 2018). In this section we augment \mathcal{B} by introducing more states, and explore its relationship to some neural architectures motivated by *n*-gram features. We start with a three-state WFSA as an example, and then discuss more general cases.

Figure 2 diagrams a WFSA C, augmenting \mathcal{B} with another state. To reach the final state q_2 , at

 $^{^{4}}$ We restrict that both operations take constant time and space, to exclude the use of arbitrarily complex semirings (§4.3).

⁵The SRU architecture discussed through this work is based on Lei et al. (2017b). In a later updated version, Lei et al. (2018) introduce diagonal matrix multiplication interaction in the hidden state updates, inspired by (Li et al., 2018), which yields a recurrence not obviously rational.



Figure 2: A three-state WFSA C discussed in §4.2.

least two transitions must be taken, in contrast to one in \mathcal{B} . History information is decayed by the self-loop at the final state q_2 , assuming ϕ_2 is between 0 and 1. \mathcal{C} has another self-loop over q_1 , weighted by $\phi_1 \in (0, 1)$. The motivation is to allow (but down-weight) nonconsecutive bigrams, as we will soon show.

The scores assigned by \mathcal{C} can be inductively computed by applying the Forward algorithm (§2). Given input sequence **x** longer than one, let $\mathcal{C}[\![\mathbf{x}_{:0}]\!] = 0$, then $\mathcal{C}[\![\mathbf{x}_{:t+1}]\!] =$

$$\mathbb{C}[\![\mathbf{x}_{:t}]\!] \phi_2(x_{t+1}) + \beta_t \,\mu_2(x_{t+1}), \qquad (11)$$

where

$$\beta_t = \beta_{t-1} \phi_1(x_t) + \mu_1(x_t), \qquad (12)$$

and $\beta_0 = 0$. Unrolling β_t in time, we get $\beta_t =$

$$\sum_{j=1}^{t-1} \left(\mu_1(x_j) \prod_{k=j+1}^t \phi_1(x_k) \right) + \mu_1(x_t). \quad (13)$$

Due to the self-loop over state q_1 , β_t can be seen as a weighted sum of the μ_1 terms up to x_t (Equaltion 13). The second product term in Equation 11 then provides multiplicative interactions between μ_2 , and the weighted sum of μ_1 s. In this sense, it captures nonconsecutive bigram features.

At a first glance, Equations 11 and 12 resemble recurrent convolutional neural networks (RCNN; Lei et al., 2016). RCNN is inspired by nonconsecutive n-gram features and low rank tensor factorization. It is later studied from a string kernel perspective (Lei et al., 2017a). Here we review its nonlinear bigram version:

$$\mathbf{c}_t^{(1)} = \mathbf{c}_{t-1}^{(1)} \odot \boldsymbol{\lambda}_t + \mathbf{u}_t^{(1)}, \qquad (14a)$$

$$\mathbf{c}_t^{(2)} = \mathbf{c}_{t-1}^{(2)} \odot \boldsymbol{\lambda}_t + \mathbf{c}_{t-1}^{(1)} \odot \mathbf{u}_t^{(2)}, \qquad (14b)$$

where the $\mathbf{u}_t^{(j)}$ s are computed similarly to Equation 5b, and $\mathbf{c}_t^{(2)}$ is used as output for onward computation. Different strategies to computing λ_t were explored (Lei et al., 2015, 2016). When λ_t is a constant, or depends only on x_t , e.g., $\lambda_t = \sigma(\mathbf{W}_{\lambda}\mathbf{v}_t + \mathbf{b}_{\lambda})$, the *i*th dimension of Equations 14



Figure 3: WFSA \mathcal{D}_1 discussed in §4.3. Two initial states q_1 and q_4 are used here.

can be recovered from Equation 11, by letting

$$\mu_j(x_t) = [\mathbf{u}_t^{(j)}]_i, \, \phi_j(x_t) = [\boldsymbol{\lambda}_t]_i, \, j = 1, 2.$$
 (15)

It is straightforward to generalize the above discussion to higher order cases: n-gram RCNN corresponds to WFSAs with n + 1 states, constructed similarly to how we build \mathcal{C} from \mathcal{B} (Appendix B).

Proposition 12. For a single-layer RCNN with λ_t being a constant or depending only on x_t , the recurrence is rational.

As noted later in §4.3, its recurrence may not be rational when $\lambda_t = \sigma(\mathbf{W}_c \mathbf{c}_{t-1} + \mathbf{W}_\lambda \mathbf{v}_t + \mathbf{b}_\lambda)$.

4.3 Beyond Elementwise Operations

So far we have discussed rational recurrences for models using elementwise recurrent updates (e.g., Equation 5c). This section uses an existing model as an example, to study a rational recurrence that is not elementwise. We focus on the input switched affine network (ISAN; Foerster et al., 2017). Aiming for efficiency and interpretability, it does not use any explicit nonlinearity; its affine transformation parameters depend only on the input:

$$\mathbf{c}_t = \mathbf{W}_{x_t} \mathbf{c}_{t-1} + \mathbf{b}_{x_t}.$$
 (16)

Due to the matrix multiplication, the recurrence of a single-layer ISAN is not elementwise. Yet, we argue that it is rational. We will sketch the proof for a 2-dimensional case, and it is straightforward to generalize to higher dimensions (Appendix C).

We define two WFSAs, each recovering one dimension of ISAN's recurrent updates. Figure 3 diagrams one of them, \mathcal{D}_1 . The other one, \mathcal{D}_2 , is identical (including shared weights), except using q_3 instead of q_2 as the final state. For any nonempty input sequence $\mathbf{x} \in \Sigma^+$, the scores assigned by \mathcal{D}_1 and \mathcal{D}_2 can be inductively computed by applying the Forward algorithm. Letting $\mathcal{D}_1[\![\mathbf{x}_{:0}]\!] = \mathcal{D}_2[\![\mathbf{x}_{:0}]\!] = 0$, for $t \ge 1$

$$\begin{bmatrix} \mathcal{D}_1[\![\mathbf{x}_{:t}]\!]\\ \mathcal{D}_2[\![\mathbf{x}_{:t}]\!] \end{bmatrix} = \widetilde{\mathbf{W}}_{x_t} \begin{bmatrix} \mathcal{D}_1[\![\mathbf{x}_{:t-1}]\!]\\ \mathcal{D}_2[\![\mathbf{x}_{:t-1}]\!] \end{bmatrix} + \widetilde{\mathbf{b}}_{x_t}, \quad (17)$$

where

$$\widetilde{\mathbf{W}}_{x_t} = \begin{bmatrix} \mu_{1,1}(x_t) & \mu_{1,2}(x_t) \\ \mu_{2,1}(x_t) & \mu_{2,2}(x_t) \end{bmatrix},$$

$$\widetilde{\mathbf{b}}_{x_t} = \begin{bmatrix} \eta_1(x_t) \\ \eta_2(x_t) \end{bmatrix}.$$
(18)

Then Equation 16, in the case of hidden size 2, is recovered by letting $\mathbf{W}_{x_t} = \widetilde{\mathbf{W}}_{x_t}$ and $\mathbf{b}_{x_t} = \widetilde{\mathbf{b}}_{x_t}$.

Proposition 13. *The recurrence of a single-layer ISAN is rational.*

Corollary 14. For a single-layer Elman network, in the absence of any nonlinearity, the recurrence is rational.

Discussion. It is known that an Elman network can approximate any recursively computable partial function (Siegelmann and Sontag, 1995). On the other hand, in their single-layer cases, WFSAs (and thus models with rational recurrences) are restricted to rational series (Schützenberger, 1961). Therefore, we hypothesize that models like Elman networks, LSTMs, and GRUs, where the recurrences depend on previous states through affine transformations followed by nonlinearities, are not rational.

This work does not intend to propose rational recurrences as a concept general enough to include most existing RNNs. Rather, we wish to study a more constrained class of methods to better understand the connections between WFSAs and RNNs. Therefore in Definition 8, we restrict the semirings to be "simple," in the sense that both operations take constant time and space. Such a restriction aims to exclude the possibility of hiding arbitrarily complex computations inside the semiring, which might allow RNNs to satisfy the definition in a trivial and unilluminating way.

Such theoretical limitations might be less severe than they appear, since it is not yet entirely clear what they correspond to in practice, especially when multiple vertical layers of these models are used (Leshno and Schocken, 1993). We defer to future work the further study of the connections between WFSAs and Elman-style RNNs.

Closing this section, Table 1 summarizes the discussed recurrent neural architectures and their corresponding WFSAs.

5 Deriving Neural Models from WFSAs

Rational recurrences provide a new view of several recently proposed neural models. Based on such

Models	Recurrence Function	WFSA
§4.1 SRU, SCRN T-RNN, QRN	$\mathbf{n} \mathbf{c}_t = \mathbf{f}_t \odot \mathbf{c}_{t-1} + \mathbf{u}_t$	В
§4.2 RCNN	$\mathbf{c}_t^{(1)} = \mathbf{c}_{t-1}^{(1)} \odot \boldsymbol{\lambda}_t + \mathbf{u}_t^{(1)}$ $\mathbf{c}_t^{(2)} = \mathbf{c}_{t-1}^{(2)} \odot \boldsymbol{\lambda}_t + \mathbf{c}_{t-1}^{(1)} \odot$	$\mathbf{u}_{t}^{(2)}$ \mathcal{C}
§4.3 ISAN	$\mathbf{c}_t = \mathbf{W}_{x_t} \mathbf{c}_{t-1} + \mathbf{b}_{x_t}$	$\mathfrak{D}_1,\mathfrak{D}_2$

Table 1: Recurrent neural network architectures discussed in §4 and their corresponding WFSAs. §4.1: SRU (Lei et al., 2017b), SCRN (Mikolov et al., 2014), T-RNN and its gated variants (Balduzzi and Ghifary, 2016), and QRNN (Bradbury et al., 2017); §4.2: RCNN (Lei et al., 2016); §4.3: ISAN (Foerster et al., 2017).

observations, this section aims to explore potential approaches to designing neural architectures in a more interpretable and intuitive way: by deriving them from WFSAs. §5.1 studies an interpolation of unigram and bigram features by combining 2state and 3-state WFSAs (Figures 1 and 2). We then explore alternative semirings (§5.2), an approach orthogonal to what we've discussed so far.

We note that our goal is not to devise new stateof-the-art architectures. Rather, we illustrate a new design process for neural architectures that draws inspiration from WFSAs. That said, in our experiments (\S 6), one of our new architectures performs as well as or better than strong baselines.

5.1 Aggregating Different Length Patterns

We start by presenting a straightforward extension to 2-state and 3-state rational models: one combining both. It is inspired by many classical NLP models, where unigram features and higher-order ones are interpolated.

Figure 4 diagrams a 4-state WFSA \mathcal{F} . Compared to \mathcal{C} (Figure 2), \mathcal{F} uses q_1 as a second final state, aiming to capture both unigram and bigram patterns, since a path is allowed to stop at q_1 after consuming one input. The final states are weighted by ρ_1 and ρ_2 respectively. Another notable modification is the additional state q_3 , which is used to create a "shortcut" to reach q_2 , together with an ε -transition. Specifically, starting from q_0 , a path can now take the ε -transition and reach q_3 , and then take a transition with weight μ_2 to reach q_2 . Recall from §2, that ε -transitions do not consume any input, yet they can still be weighted by a (parameterized) function γ not depending on the inputs. The ε -transition allows for skipping the



Figure 4: A WFSA \mathcal{F} that combines both unigram and bigram features (§5.1). Two final states q_1 and q_2 are used, with weights ρ_1 and ρ_2 , respectively.

first word in a bigram. It can be discouraged by using $\gamma \in (0, 1)$, just as we do in our experiments.

Deriving the neural architecture. As in §3, we relate hidden states of an RNN to the scores assigned by WFSAs to input strings. We then derive the neural architecture with a dynamic program. Here we keep the discussion self-contained by explicitly overviewing the procedure. It is a direct application of the Forward algorithm (§2), though now in a form that deals with the ε -transition. Such an approach applies, of course, to more general cases, as noted by Schwartz et al. (2018).

Given an input string $\mathbf{x} \in \Sigma^+$, let $z_t^{(j)}$ denote the total score of all paths landing in state q_j just after consuming x_t . Let $z_0^{(j)} = 0$, then for $t \ge 1$,

$$\begin{aligned} z_t^{(0)} &= 1\\ z_t^{(1)} &= z_{t-1}^{(1)} \phi_1(x_t) + z_{t-1}^{(0)} \mu_1(x_t)\\ z_t^{(3)} &= z_t^{(0)} \gamma\\ z_t^{(2)} &= z_{t-1}^{(2)} \phi_2(x_t) + (z_{t-1}^{(1)} + z_{t-1}^{(3)}) \, \mu_2(x_t)\\ \mathcal{F}[\![\mathbf{x}_{:t}]\!] &= \rho_1 \, z_t^{(1)} + \rho_2 \, z_t^{(2)}. \end{aligned}$$

We now collect d of these WFSAs to construct an RNN, and we parameterize their weight functions with the technique we've been using:

$$\mathbf{c}_{t}^{(1)} = \mathbf{c}_{t-1}^{(1)} \odot \mathbf{f}_{t}^{(1)} + \mathbf{u}_{t}^{(1)}, \tag{19a}$$

$$\mathbf{c}_{t}^{(2)} = \mathbf{c}_{t-1}^{(2)} \odot \mathbf{f}_{t}^{(2)} + (\mathbf{c}_{t-1}^{(1)} + \mathbf{r}) \odot \mathbf{u}_{t}^{(2)}, \quad (19b)$$

$$\mathbf{c}_t = \mathbf{p}^{(1)} \odot \mathbf{c}_t^{(1)} + \mathbf{p}^{(2)} \odot \mathbf{c}_t^{(2)}, \qquad (19c)$$

where

$$\mathbf{f}_t^{(j)} = \boldsymbol{\sigma} \big(\mathbf{W}_f^{(j)} \mathbf{v}_t + \mathbf{b}_f^{(j)} \big), \tag{20a}$$

$$\mathbf{u}_t^{(j)} = (\mathbf{1} - \mathbf{f}_t^{(j)}) \odot \boldsymbol{g} \big(\mathbf{W}_u^{(j)} \mathbf{v}_t + \mathbf{b}_u^{(j)} \big), \quad (20b)$$

$$\mathbf{p}^{(j)} = \boldsymbol{\sigma}(\mathbf{b}_p^{(j)}), \quad \mathbf{r} = \boldsymbol{\sigma}(\mathbf{b}_r).$$
 (20c)

The p vectors correspond to the final state weights ρ_1 and ρ_2 . Despite the similarities, p are different from output gates (Bradbury et al., 2017), since the former do not depend on the input, and are param-

Model	Unigram	Bigram	Semiring
rrnn(B)	\checkmark		real
$\operatorname{RRNN}(\mathcal{B})_{m+}$	\checkmark		max-plus
rrnn(C)		\checkmark	real
$\operatorname{RRNN}(\mathcal{F})$	\checkmark	\checkmark	real

Table 2: Rational recurrent neural architectures compared in the experiments $(\S 6.1)$.

eterized (through a sigmoid) by two leaned vectors $\mathbf{b}_p^{(j)}$. The same applies to \mathbf{r} and \mathbf{b}_r , which correspond to the weights for ε -transitions γ .

5.2 Alternative Semirings

Our new understanding of rational recurrences allows us to consider a different kind of extension: replacing the semiring. We introduce an example, which modifies Example 6 by replacing its real (plus-times) semiring with the max-plus semiring $\langle \mathbb{R} \cup \{-\infty\}, \max, +, -\infty, 0 \rangle$:

Example 15.

$$\mathbf{f}_t = \log \boldsymbol{\sigma} (\mathbf{W}_f \mathbf{v}_t + \mathbf{b}_f),$$
 (21a)

$$\mathbf{u}_t = \boldsymbol{g} \big(\mathbf{W}_u \mathbf{v}_t + \mathbf{b}_u \big),$$
 (21b)

$$\mathbf{c}_t = \max\{\mathbf{f}_t + \mathbf{c}_{t-1}, \mathbf{u}_t\}.$$
 (21c)

Example 15 does not use the forget gate when computing \mathbf{u}_t (Equation 21b), which is different from its plus-times counterpart, where $\mathbf{u}_t =$ $(\mathbf{1} - \mathbf{f}_t) \odot g(\mathbf{W}_u \mathbf{v}_t + \mathbf{b}_u)$. The reason is that, unlike the real semiring, the max-plus semiring lacks a well-defined negation. Possible alternatives include taking the log of a separate input gate, or using $\log(\mathbf{1} - \mathbf{f}_t)$, which we leave for future work.

Example 15 can be seen as replacing sumpooling with max-pooling. Both max and sumpooling have been used successfully in vision and NLP models. Intuitively, max-pooling "detects" the occurrence of a pattern while sum-pooling "counts" the occurrence of a pattern. One advantage of max operator is that the model's decisions can be back-traced and interpreted, as argued by Schwartz et al. (2018). Such a technique is applicable to all the models with rational recurrences.

6 Experiments

This section evaluates four rational RNNs on language modeling ($\S6.2$) and text categorization ($\S6.3$). Our goal is to compare the behaviors of models derived from different WFSAs, showing that our understanding of WFSAs allows us to

improve existing rational models.

6.1 Compared Models

Our comparisons focus on the *recurrences* of the models, i.e., how the hidden states \mathbf{c}_t are computed (e.g., Equations 5c and 19c). Therefore we follow Lei et al. (2017b) and use $\mathbf{u}_t^{(j)} = \mathbf{W}_u^{(j)} \mathbf{v}_t^{(j)}$ across all compared models, listed below and as well as in Table 2:

- RRNN(\mathcal{B}), with real semiring (§4.1);
- RRNN(\mathcal{B})_{m+}, with max-plus semiring (§5.2);
- RRNN(C), with real semiring (§4.2);
- RRNN(\mathcal{F}), with real semiring (§5.1).

We also compare to an LSTM baseline. Aiming to control for comfounding factors, we do not use highway connections in any of the models.⁶ In the interest of space, the full architectures and hyperparameters are detailed in Appendices D and E.

6.2 Language Modeling

Dataset and implementation. We experiment with the Penn Treebank corpus (PTB; Marcus et al., 1993). We use the preprocessing and splits from Mikolov et al. (2010), resulting in a vocabulary size of 10K and 1M tokens.

Following standard practice, we treat the training data as one long sequence, split into mini batches, and train using BPTT truncated to 35 time steps (Williams and Peng, 1990). The input embeddings and output softmax weights are tied (Press and Wolf, 2017).

Results. Following Collins et al. (2017) and Melis et al. (2018), we compare models controlling for parameter budget. Table 3 summarizes language modeling perplexities on PTB test set. The middle block compares all models with two layers and 10M trainable parameters. RRNN(\mathcal{B}) and RRNN(\mathcal{C}) achieve roughly the same performance; interpolating both unigram and bigram features, RRNN(\mathcal{F}) outperforms others by more than 2.9 test perplexity. For the three-layer and 24M setting (the bottom block), we observe similar trends, except that RRNN(\mathcal{C}) slightly underperforms RRNN(\mathcal{B}). Here RRNN(\mathcal{F}) outperforms others by more than 2.1 perplexity.

Using a max-plus semiring, $RRNN(\mathcal{B})_{m+}$ underperforms $RRNN(\mathcal{B})$ under both settings. Possible reasons could be the suboptimal design choice

Model	l	# Params.	Dev.	Test
LSTM	2	24M	73.3	71.4
LSTM	3	24M	78.8	76.2
rrnn(B)	2	10M	73.1	69.2
$\operatorname{RRNN}(\mathcal{B})_{m+}$	2	10M	75.1	71.7
rrnn(C)	2	10M	72.5	69.5
$\operatorname{rrnn}(\mathcal{F})$	2	10M	69.5	66.3
rrnn(B)	3	24M	68.7	65.2
$\operatorname{RRNN}(\mathcal{B})_{m+}$	3	24M	70.8	66.9
rrnn(C)	3	24M	70.0	67.0
$\operatorname{RRNN}(\mathcal{F})$	3	24M	66.0	63.1

Table 3: Language modeling perplexity on PTB test set (lower is better). LSTM numbers are taken from Lei et al. (2017b). ℓ denotes the number of layers. Bold font indicates best performance.

Split	Amazon	SST	subj	CR
Train	20K	6.9K	8K	3.0K
Dev.	5K	0.9K	1K	0.4K
Test	25K	1.8K	1K	0.4K

Table 4: Number of instances in the text classification datasets ($\S 6.3$).

for computing input representations in the former ($\S5.2$). Finally, most compared models outperform the LSTM baselines, whose numbers are taken from Lei et al. (2017b).⁷

6.3 Text Classification

Implementation. We use unidirectional 2-layer architectures for all compared models. To build the classifiers, we feed the final RNN hidden states into a 2-layer tanh-MLP. Further implementation details are described in Appendix E.

Datasets. We experiment with four binary text classification datasets, described below.

- Amazon (electronic product review corpus; McAuley and Leskovec, 2013).⁸ We focus on the positive and negative reviews.
- **SST** (Stanford sentiment treebank; Socher et al., 2013).⁹ We focus on the binary classification task. SST provides labels for syntactic phrases; we experiment with a more realistic setup, and

⁶Thus RRNN(\mathfrak{B}) is essentially an SRU without highway connections. We denote it differently, to note its differences from the original implementation (Lei et al., 2017b). Similarly, we do not denote RRNN(\mathfrak{C}) as RCNN (Lei et al., 2016).

⁷Melis et al. (2018) point out that carefully tuning LSTMs can achieve much stronger performance, at the cost of exceptionally large amounts of computational resources for tuning. ⁸http://riejohnson.com/cnn_data.html

⁹nlp.stanford.edu/sentiment/index.html

Model	Amazon	SST	subj	CR
LSTM	$91.2_{\pm0.3}$	$85.1_{\pm 0.6}$	$93.3_{\pm 0.6}$	$82.4_{\pm 1.5}$
rrnn(多) rrnn(C)			$93.9_{\pm 0.4}\\93.8_{\pm 0.6}$	
rrnn(B) _{m+} rrnn(牙)	$\begin{array}{c} 89.2_{\pm 3.1} \\ 92.7_{\pm 0.2} \end{array}$			

Table 5: Text classification test accuracy averaged over 5 runs. \pm denotes standard deviation, and bold font indicates best averaged performance.

consider only complete sentences at either training or evaluating time.

- **subj** (Subjectivity dataset; Pang and Lee, 2004). As **subj** doesn't come with official splits, we randomly split it to train (80%), development (10%), and test (10%) sets.
- **CR** (customer reviews dataset; Hu and Liu, 2004).¹⁰ As with **subj**, we randomly split this dataset using the same ratio.

Table 4 summarizes the sizes of the datasets.

Results. Table 5 summarizes text classification test accuracy. We report the average performance of 5 trials different only in random seeds. $RRNN(\mathcal{F})$ outperforms all other models on 3 out of the 4 datasets. For Amazon, the largest one, we do not observe significant differences between $RRNN(\mathcal{F})$ and $RRNN(\mathcal{C})$, while both outperform others. This may suggest that the interpolation of unigram and bigram features by $RRNN(\mathcal{F})$ is especially useful in small data setups. As in the language modeling experiments, $RRNN(\mathcal{B})_{m+}$ underperforms all other models in most cases, and in particular RRNN(B). These results provide evidence that replacing the real semiring in rational models might be challenging. We leave further exploration to future work.

7 Related Work

Weighted finite state automata. WFSAs were once popular among many sequential tasks (Mohri et al., 2002; Kumar and Byrne, 2003; Cortes et al., 2004; Pardo and Birmingham, 2005; Moore et al., 2006, *inter alia*), and are still successful in morphology (Dreyer, 2011; Cotterell et al., 2015; Rastogi et al., 2016, *inter alia*). Compared to neural networks, WFSAs are better understood theoretically and arguably more interpretable. They were recently revisited in combination with the former in, e.g., text generation (Ghazvininejad et al., 2016, 2017; Lin et al., 2017) and automatic music accompaniment (Forsyth, 2016).

Recurrent neural networks. RNNs (Elman, 1990; Jordan, 1989) prove to be strong models for sequential data (Siegelmann and Sontag, 1995). Besides the perhaps most notable gated variants (Hochreiter and Schmidhuber, 1997; Cho et al., 2014), extensive efforts have been devoted to developing alternatives (Balduzzi and Ghifary, 2016; Miao et al., 2016; Zoph and Le, 2017; Lee et al., 2017; Lei et al., 2017a; Vaswani et al., 2017; Gehring et al., 2017, *inter alia*). Departing from the above approaches, this work derives RNN architectures drawing inspiration from WFSAs.

Another line of work studied the connections between WFSAs and RNNs in terms of modeling capacity, both empirically (Kolen, 1993; Giles et al., 1992; Weiss et al., 2018, *inter alia*) and theoretically (Cleeremans et al., 1989; Visser et al., 2001; Chen et al., 2018, *inter alia*).

8 Conclusion

We presented *rational recurrences*, a new construction to study the recurrent updates in RNNs, drawing inspiration from WFSAs. We showed that rational recurrences are in frequent use by several recently proposed recurrent neural architectures, providing new understanding of them. Based on such connections, we discussed approaches to deriving novel neural architectures from WFSAs. Our empirical results demonstrate the potential of doing so. We publicly release our implementation at https://github.com/ Noahs-ARK/rational-recurrences.

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¹⁰http://www.cs.uic.edu/?liub/FBS/ sentiment-analysis.html

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