

### On the Practical Computational Power of Finite Precision RNNs for Language Recognition

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### **GRU < LSTM (!?)**



Supported by European Union's Seventh Framework Programme (FP7) under grant agreement no. 615688 (PRIME)

### **Current State**

- RNNs are everywhere
- We don't know too much about the differences between them:
  - Gated RNNs are shown to train better, beyond that:
  - "RNNs are Turing Complete"?

# **Turing Complete?**

#### On the Computational Power of Neural Nets\*

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Received February 4, 1992; revised May 24, 1993

## **Turing Complete?**

1993 Proof:

#### **1. Requires Infinite Precision:**

**Uses stack(s), maintained in certain dimension(s) Zeros are pushed using division** (using g = g/4 + 1/4) In 32 bits, this reaches the limit after **15** pushes

#### **2. Requires Infinite Time:**

Allows processing steps beyond reading input (Not the standard use case!)

unreasonable assumptions!

## **Turing Complete?**



unreasonable assumptions!

What happens on real hardware and real use-cases?

### Real Use

- Gated architectures have the best performance
  - LSTM and GRU are most popular
  - Of these, the choice between them is unclear

### Main Result

#### We accept all RNN types can simulate DFAs

#### We show that LSTMs and IRNNs can also count

#### And that the GRU and SRNN cannot

### Power of Counting

**Practical** 

#### In NMT: LSTM better at capturing target length

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#### Theoretical

Finite State Machines vs Counter Machines

#### K-Counter Machines (SKCMs)

Fischer, Meyer, Rosenberg - 1968

- Similar to finite automata, but also maintain k counters
- A counter has 4 operations: inc/dec by one, do nothing, reset
- Counters are observed by comparison to zero



### Counting Machines and Chomsky Hierarchy











#### **SKCMs cross the Chomsky Hierarchy!**



## Summary so Far

- Counters give additional formal power
- We claimed that LSTM can count and GRU cannot

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- Counters give additional formal power
- We claimed that LSTM can count and GRU cannot
- Let's see why

$$z_{t} = \sigma(W^{z}x_{t} + U^{z}h_{t-1} + b^{z})$$

$$r_{t} = \sigma(W^{r}x_{t} + U^{r}h_{t-1} + b^{r})$$

$$\tilde{h}_{t} = \tanh(W^{h}x_{t} + U^{h}(r_{t} \circ h_{t-1}) + b^{h})$$

$$h_{t} = z_{t} \circ h_{t-1} + (1 - z_{t}) \circ \tilde{h}_{t}$$

$$f_t = \sigma(W^f x_t + U^f h_{t-1} + b^f)$$
  

$$i_t = \sigma(W^i x_t + U^i h_{t-1} + b^i)$$
  

$$o_t = \sigma(W^o x_t + U^o h_{t-1} + b^o)$$
  

$$\tilde{c}_t = \tanh(W^c x_t + U^c h_{t-1} + b^c)$$
  

$$c_t = f_t \circ c_{t-1} + i_t \circ \tilde{c}_t$$
  

$$h_t = o_t \circ g(c_t)$$

$$\begin{aligned} z_t &= \sigma(W^z x_t + U^z h_{t-1} + b^z) \\ r_t &= \sigma(W^r x_t + U^r h_{t-1} + b^r) \end{aligned} \qquad \text{gates} \qquad f_t &= \sigma(W^J x_t + U^J h_{t-1} + b^J) \\ \tilde{h}_t &= tanh(W^h x_t + U^h(r_t \circ h_{t-1}) + b^h) \\ \tilde{h}_t &= tanh(W^h x_t + U^h(r_t \circ h_{t-1}) + b^h) \\ h_t &= z_t \circ h_{t-1} + (1 - z_t) \circ \tilde{h}_t \end{aligned} \qquad \tilde{h}_t \qquad \tilde{c}_t &= tanh(W^c x_t + U^c h_{t-1} + b^c) \\ \tilde{c}_t &= tanh(W^c x_t + U^c h_{t-1} + b^c) \\ \tilde{c}_t &= f_t \circ c_{t-1} + i_t \circ \tilde{c}_t \\ h_t &= o_t \circ g(c_t) \end{aligned}$$



→ gates →  $f_t \in (0,1)$  $z_t \in (0,1)$  $i_t \in (0,1)$  $r_t \in (0,1)$  $o_t \in (0,1)$  $\tilde{h}_t \in (-1,1)$  $\tilde{c}_t \in (-1,1)$  $h_t = z_t \circ h_{t-1} + (1 - z_t) \circ h_t$  $c_t = f_t \circ c_{t-1} + i_t \circ \tilde{c}_t$ candidate  $h_t = o_t \circ g(c_t)$ vectors update functions

 $z_t \in (0,1)$   $r_t \in (0,1)$   $\tilde{h}_t \in (-1,1)$  $h_t = z_t \circ h_{t-1} + (1-z) \circ \tilde{h}_t$   $f_{t} \in (0,1)$   $i_{t} \in (0,1)$   $o_{t} \in (0,1)$   $\tilde{c}_{t} \in (-1,1)$   $c_{t} = f_{t} \circ c_{t-1} + i_{t} \circ \tilde{c}_{t}$   $h_{t} = o_{t} \circ g(c_{t})$ 

 $z_{t} \in (0,1)$   $r_{t} \in (0,1)$   $\tilde{h}_{t} \in (-1,1)$  $h_{t} = z_{t} \circ h_{t-1} + (1-z) \circ \tilde{h}_{t}$   $f_{t} \in (0,1)$   $i_{t} \in (0,1)$   $o_{t} \in (0,1)$   $\tilde{c}_{t} \in (-1,1)$   $c_{t} = f_{t} \circ c_{t-1} + i_{t} \circ \tilde{c}_{t}$  $h_{t} = o_{t} \circ g(c_{t})$ 

#### Interpolation

 $z_{t} \in (0,1)$   $r_{t} \in \textbf{Boundedl}$   $\tilde{h}_{t} \in (-1,1)$  $h_{t} = z_{t} \circ h_{t-1} + (1-z) \circ \tilde{h}_{t}$ 

#### Interpolation

 $f_{t} \in (0,1)$   $i_{t} \in (0,1)$   $o_{t} \in (0,1)$   $\tilde{c}_{t} \in (-1,1)$   $c_{t} = f_{t} \circ c_{t-1} + i_{t} \circ \tilde{c}_{t}$   $h_{t} = o_{t} \circ g(c_{t})$ 

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Interpolation

 $f_{t} \in (0,1)$   $i_{t} \in (0,1)$   $o_{t} \in (0,1)$   $\tilde{c}_{t} \in (-1,1)$   $c_{t} = f_{t} \circ c_{t-1} + i_{t} \circ \tilde{c}_{t}$  $h_{t} = o_{t} \circ g(c_{t})$ 

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#### Interpolation

 $f_{t} \in (0,1)$   $i_{t} \in (0,1)$   $o_{t} \in (0,1)$   $\tilde{c}_{t} \in (-1,1)$   $c_{t} = f_{t} \circ c_{t-1} + i_{t} \circ \tilde{c}_{t}$  $h_{t} = o_{t} \circ g(c_{t})$ 

#### Addition

 $z_{t} \in (0,1)$   $r_{t} \in \textbf{Bounded!}$   $\tilde{h}_{t} \in (-1,1)$   $h_{t} = z_{t} \circ h_{t-1} + (1-z) \circ \tilde{h}_{t}$ 

#### Interpolation

 $f_t \approx 1$   $i_t \approx 1$   $o_t \in (0,1)$   $\tilde{c}_t \in (-1,1)$   $c_t \approx c_{t-1} + \tilde{c}_t$  $h_t = o_t \circ g(c_t)$ 

Addition

$$c_t = f_t \circ c_{t-1} + i_t \circ \tilde{c}_t$$

 $z_{t} \in (0,1)$   $r_{t} \in \textbf{Bounded!}$   $\tilde{h}_{t} \in (-1,1)$   $h_{t} = z_{t} \circ h_{t-1} + (1-z) \circ \tilde{h}_{t}$ 

Interpolation

 $f_t \approx 1$   $i_t \approx 1$   $o_t \in (0,1)$   $\tilde{c}_t \approx 1$   $c_t \approx c_{t-1} + 1$  $h_t = o_t \circ g(c_t)$ 

Increase by 1

$$c_t = f_t \circ c_{t-1} + i_t \circ \tilde{c}_t$$

 $z_{t} \in (0,1)$   $r_{t} \in \textbf{Bounded!}$   $\tilde{h}_{t} \in (-1,1)$   $h_{t} = z_{t} \circ h_{t-1} + (1-z) \circ \tilde{h}_{t}$ 

Interpolation

 $f_t \approx 1$   $i_t \approx 1$   $o_t \in (0,1)$   $\tilde{c}_t \approx -1$   $c_t \approx c_{t-1} - 1$  $h_t = o_t \circ g(c_t)$ 

#### Decrease by 1

$$c_t = f_t \circ c_{t-1} + i_t \circ \tilde{c}_t$$

 $z_{t} \in (0,1)$   $r_{t} \in \textbf{Bounded!}$   $\tilde{h}_{t} \in (-1,1)$   $h_{t} = z_{t} \circ h_{t-1} + (1-z) \circ \tilde{h}_{t}$ 

#### Interpolation

 $f_t \approx 1$   $i_t \approx 0$   $o_t \in (0,1)$   $\tilde{c}_t \in (-1,1)$   $c_t \approx c_{t-1}$  $h_t = o_t \circ g(c_t)$ 

#### **Do Nothing**

$$c_t = f_t \circ c_{t-1} + i_t \circ \tilde{c}_t$$

 $z_{t} \in (0,1)$   $r_{t} \in \textbf{Bounded!}$   $\tilde{h}_{t} \in (-1,1)$   $h_{t} = z_{t} \circ h_{t-1} + (1-z) \circ \tilde{h}_{t}$ 

#### Interpolation

 $f_t \approx 0$   $i_t \approx 0$   $o_t \in (0,1)$   $\tilde{c}_t \in (-1,1)$   $c_t \approx 0$  $h_t = o_t \circ g(c_t)$ 

Reset

$$c_t = f_t \circ c_{t-1} + i_t \circ \tilde{c}_t$$

 $z_{t} \in (0,1)$   $r_{t} \in \textbf{Bounded!}$   $\tilde{h}_{t} \in (-1,1)$   $h_{t} = z_{t} \circ h_{t-1} + (1-z) \circ \tilde{h}_{t}$ 

Interpolation

 $f_t \approx 0$   $i_t \approx 0$   $o_t \in \textbf{Can Count!}$   $o_t \in (-1,1)$   $c_t \approx 0$  $h_t = o_t \circ g(c_t)$ 

Reset

$$c_t = f_t \circ c_{t-1} + i_t \circ \tilde{c}_t$$

### Other Architectures

#### SRNN

#### IRNN

 $h_t = \sigma_h(W_h x_t + U_h h_{t-1} + b_h) \qquad h_t = \max(0, W_h x_t + U_h h_{t-1} + b_h)$ 

### Other Architectures

#### SRNN

#### IRNN

 $h_t = \sigma_h(W_h x_t + U_h h_{t-1} + b_h) \in (0,1) \qquad h_t = \max(0, W_h x_t + U_h h_{t-1} + b_h)$ 


# Other Architectures

## SRNN

## IRNN



(subtraction in parallel, also increasing, counter)

## Other Architectures

## SRNN

## IRNN



(subtraction in parallel, also increasing, counter)

## So:

- LSTM can count!
- GRU cannot
- Counting gives greater computational power

**Trained**  $a^n b^n$ , (on positive examples up to length 100)

Activations on  $a^{1000}b^{1000}$  :



GRU

**Trained**  $a^n b^n$ , (on positive examples up to length 100)

Activations on  $a^{1000}b^{1000}$  :



• Took much longer to train

**Trained**  $a^n b^n$ , (on positive examples up to length 100)

Activations on  $a^{1000}b^{1000}$  :



#### 42

- Took much longer to train
- Did not generalise even within training domain
  - begin failing at n=39 (vs 257 for LSTM)

**Trained**  $a^n b^n$ , (on positive examples up to length 100)

Activations on  $a^{1000}b^{1000}$  :



#### **GRU:**

- Took much longer to train
- Did not generalise even within training domain
  - begin failing at n=39 (vs 257 for LSTM)
- Did not learn any discernible counting mechanism

**Trained**  $a^n b^n c^n$ , (on positive examples up to length 50)

Activations on  $a^{100}b^{100}c^{100}$ :



**Trained**  $a^n b^n c^n$ , (on positive examples up to length 100)

Activations on  $a^{100}b^{100}c^{100}$ :



#### **GRU**:

- Took much longer to train
- Did not generalise well
  - begin failing at n=9 (vs 101 for LSTM)
- Did not learn any discernible counting mechanism

## Conclusion

# IRNN LSTM SRNN GRU Trainability



# Take Home Message

## **Architectural Choices Matter!**

and result in actual differences in expressive power

## **Don't fall in the Turing Tarpit!**

## Thank You

## **GitHub repository:**

https://github.com/tech-srl/counting\_dimensions

### Google Colab (link through GitHub as well):

https://tinyurl.com/ybjkumrz