## A Sinhala Word Joiner

Rajith Priyanga Surangika Ranatunga Gihan Dias<br>Department of Computer Science and Engineering<br>University of Moratuwa<br>Sri Lanka<br>rpriyanga@yahoo.com, surangika@cse.mrt.ac.lk, gihan@uom.lk


#### Abstract

Sinhala is an agglutinative language where many words are formed by joining several morphemes. Word joining is a basic operation in Sinhala morphology, and is based on sandhi rules.


The Sinhala word joiner is a software component which implements sandhi rules to synthesise a word from two or more morphemes. We studied Sinhala word join rules based on grammar and usage and implemented a library and a standalone application to synthesise Sinhala words. In addition to the joined word, it also outputs the rule used for joining. The tool uses a combination of a corpus and hand-coded rules to improve accuracy.

Keywords: Sinhala word joiner, Morphophonemic tools, corpus based scoring algorithm, sandhi rules

## 1 Introduction

Sinhala belongs to the Indo Aryan sub branch of the Indo-European language family. It is a descendent of the Sanskrit language, but was heavily influenced by the Pāli language from the second century B.C. as a result of the introduction of Buddhism to Sri Lanka. Other than from Pāli, Sinhala was influenced mainly by Tamil, Arabic, Portuguese, Dutch and English languages. Sinhala is written in its own script which is a descendent of the Brahmi script.

Even though the Sinhala language and its script have many similarities with their ancestors from India, they have evolved uniquely over two millennia.

There have been some attempts to implement morphological synthesizers and analysers fЪßO
minis is the lemma of the noun man．$u$ is the suffix to generate the plural subject form of the noun．

In derivational morphology，words of different word classes or with different meanings are formed by applying morphological operations on the lemma．

duk means suffering．pat means become．The combined word duppat is an adjective that means poor．

Based on the two－level morphology concept， these morphological operations can be represented in two stages．（Kimmo，1984）

1．Lexical Representation
2．Surface Representation e．g．：
1．Lexical representation of the plural subject form of the noun lemma minis is（minis $+u$ ） （ ఠోో + C）

2．Surface representation of（minis $+u$ ）is minissu（ฮెోజోజ్ష్ ）

For morphological synthesis，both lexical and surface representation rules should be applied． Surface representation rules are generally based on phonology and may transform both the left－ and right－hand morphemes．This transformation is called morphophonemics．

In Sinhala，the most common types of word joining are：
prefix + word $\rightarrow$ word
word（or lemma）＋suffix $\rightarrow$ word
word + word $\rightarrow$ word
Where + is the join operator．
Most inflectional morphology operations are of the＂lemma + suffix $\rightarrow$ word＂form．

Many derivational morphology operations are of the＂prefix＋word $\rightarrow$ word＂and＂word＋word $\rightarrow$ word＂form．


Meaning is as explained above．

๑งエ゙とがか）
pol means coconut．atta means branch．pollatta means the branch of a coconut tree．

In Sinhala，this set of morphophonemic rules are called sandhi（join rules）．Similar to the Ashṭādhyāȳ̄ of Pānuini in Sanskrit，The Sinhala grammar book Sidat San̆garā written in 13th century A．D．by Vēdēha Swāmi describes some of the grammatical aspects of the Sinhala language．

There are nine join rules in Sinhala language according to Sidat San̆garā．

Following is an example of how the Sidat San̆gara $\bar{a}$ has explained join rules．This join rule is named Pūrwa Swạra Lōpạ．

## ＂pera sara lopā para sara gatạta pceminạ＂


Meaning：Vowel part of the last letter of the left word is replaced by the first vowel in the second word．
According to the above definition，there are two conditions for this rule to be valid．

1．The last letter of the left word must be a combined letter that has a consonant and a vowel part
2．The first letter of the right word must be a vowel．
The other join rules are similarly defined．

## 2．1 Sinhala Word Join Rules

We represent the join rules described in Sidat Saňgarā in an easily understandable format as follows．
Where
－ $\mathrm{Ci}=$ consonant（e．g．$\left.k-\infty^{p}\right)$
－ $\mathrm{Vi}=$ vowels（e．g．$a-$ व̈）
－Individual letters at the word boundary are written in square brackets．（e．g．［C1］）
－Combined letters that have a consonant and a vowel in it is written in $[\mathrm{Ci} \mid \mathrm{Vi}]$ form．
e．g．：$k a=[k \mid a](\infty=\infty+q)$
－$L$ and $R$ are the remaining parts of the joining morphemes．

## 1. Pūrwạ Swạrạ Lōpạ

$\mathrm{L}[\mathrm{C} 1 \mid \mathrm{V} 1]+[\mathrm{V} 2] \mathrm{R} \rightarrow \mathrm{L}[\mathrm{C} 1 \mid \mathrm{V} 2] \mathrm{R}$

## 2. Parạ Swạrạ Lōpạ <br> $\mathrm{L}[\mathrm{C} 1 \mid \mathrm{V} 1]+[\mathrm{V} 2] \mathrm{R} \rightarrow \mathrm{L}[\mathrm{C} 1 \mid \mathrm{V} 1] \mathrm{R}$

## 3. Swạa

$\mathrm{L}[\mathrm{C} 1]+[\mathrm{V} 1] \mathrm{R} \rightarrow \mathrm{L}[\mathrm{C} 1 \mid \mathrm{V} 1] \mathrm{R}$

## 4. Swạrādeshạ

$\mathrm{L}[\mathrm{C} 1 \mid a]+[i] \mathrm{R} \rightarrow \mathrm{L}[\mathrm{C} 1 \mid e] \mathrm{R}$
$\mathrm{L}[\mathrm{C} 1 \mid a]+[u] \mathrm{R} \rightarrow \mathrm{L}[\mathrm{C} 1 \mid o] \mathrm{R}$
$\mathrm{L}[\mathrm{C} 1 \mid a]+[u] \mathrm{R} \rightarrow \mathrm{L}[\mathrm{C} 1 \mid \bar{o}] \mathrm{R}$

## 5. Gatrādeshạ

$\mathrm{L}[\mathrm{C} 1 \mid \mathrm{V} 1]+[\mathrm{C} 2 \mid \mathrm{V} 2] \mathrm{R} \rightarrow \mathrm{L}[\mathrm{C} 1 \mid \mathrm{V} 1][\mathrm{C} 3 \mid \mathrm{V} 2] \mathrm{R}$
Where C3 is a member of $\{y, v, h, k, t, p, n$, $m\}$
6. Pūrwạ Rūpa
$\mathrm{L}[\mathrm{C} 1]+[\mathrm{C} 2 \mid \mathrm{V} 2] \mathrm{R} \rightarrow \mathrm{L}[\mathrm{C} 1][\mathrm{C} 1 \mid \mathrm{V} 2] \mathrm{R}$
7. Gatrākshạrạ Lōpạ
$\mathrm{L}[n]+[\mathrm{C} 2 \mid \mathrm{V} 2] \mathrm{R} \rightarrow \mathrm{L}[\check{n} g \mid \mathrm{V} 2] \mathrm{R}$
$\mathrm{L}[n]+[\mathrm{C} 2 \mid \mathrm{V} 2] \mathrm{R} \rightarrow \mathrm{L}[\check{n} b \mid \mathrm{V} 2] \mathrm{R}$
8. Āgạma
$\mathrm{L}[\mathrm{C} 1]+\mathrm{R} \rightarrow \mathrm{L}[\mathrm{C} 1 \mid \mathrm{u}] \mathrm{R}$
$\mathrm{L}[\mathrm{C} 1]+\mathrm{R} \rightarrow \mathrm{L}[\mathrm{C} 1 \mid \mathrm{i}] \mathrm{R}$
$\mathrm{L}[\mathrm{C} 1 \mid \mathrm{V} 1]+[\mathrm{V} 2] \mathrm{R} \rightarrow \mathrm{L}[\mathrm{C} 1 \mid \mathrm{V} 1][\mathrm{C} 3 \mid \mathrm{V} 2] \mathrm{R}$
Where C3 $=\{y, v, r\}$

## 9. Dvitvạ Rūpạ

$$
\mathrm{L}[\mathrm{C} 1 \mid \mathrm{V} 1]+[\mathrm{V} 2] \mathrm{R} \rightarrow \mathrm{~L}[\mathrm{C} 1][\mathrm{C} 1 \mid \mathrm{V} 2] \mathrm{R}
$$

In addition to the above 9 join rules, we identified a few more join rules in current Sinhala. Some of them are directly taken from Sanskrit and are used in loanwords. The following join rule is an example of a rule that is not in Sidat Saňgarā, but currently in use. (Karunathilaka, 1995)

## 11. Parạ Rūpa

$\mathrm{L}[\mathrm{C} 1]+[\mathrm{C} 2 \mid \mathrm{V} 2] \mathrm{R} \rightarrow \mathrm{L}[\mathrm{C} 2][\mathrm{C} 2 \mid \mathrm{V} 2] \mathrm{R}$

## 3 Previous Work

In implementing an English to Sinhala machine translator, Hettige and Karunananda (2011) have implemented a morphological synthesizer. They generate all the forms of all noun classes considering the changes to the letters at the word boundaries in the transformation. They have not used generic joining rules for joining Sinhala words and morphemes but have defined a large number of specific finite state automata to handle multiple letter combination at the word boundaries. However, they do not cover all combinations.

To obtain the indistinct singular subject form of the noun lemma miti (short person) and balu (dog) they implement 2 different automata, which result in mittā and ballā respectively.

Remove $t i$ (రి) and append $t t \bar{a}$ (১อృ)

Remove $l u\left(\mathrm{e}_{\mathrm{a}}\right)$ and append $l l \bar{a}$ (e) $\mathrm{C}^{\mathrm{C}}$ )
They have implemented 85 FSA for Sinhala noun formations. However, both of the above transformations use the common join rule called Dvitva Rūpa, and may be defined as a single FSA.

Also in their method, the FSA must be input to the noun form synthesizer. For the same letter combinations at word boundaries, different finite state automata must be used for different word morpheme combinations. It is not possible to locate the correct FSA without a comprehensive knowledge of the Sinhala language.

There are no other significant work done in the area of Sinhala morphological synthesis or word joining.

Word joiners have been implemented for other Indic languages such as Hindi and Sanskrit. (Jha et al., 2009; Hyman, 2009; Gupta and Goyal, 2017; Kumar et al., 2010) Some of them use the join rules mentioned in the $A s h t \bar{a} d h y \bar{a} y \bar{\imath}$ of Pāṇini. (Jha et al., 2009; Hyman, 2009; Gupta and Goyal, 2017)

Most of them have used finite state transducers to do the morphophonemic operations to obtain the surface form. Most of the morphological tool
implementations for European languages also use finite state transducers to obtain the surface representation for the lexical representation (Lauri et al., 1992). Finite state transducers have been widely used in solving this morphology problem in different language families.

## 4 Challenges

In Sinhala, for a given pair of morphemes, there may be multiple matching join rules based on the boundary conditions. Also, even when a single join rule is applied, there can be multiple possible outputs, all of which are not necessarily valid for a given pair.

Accordingly, we have identified the following scenarios for a pair of or morphemes.

1. There is only one matching join rule for the pair. The combined form/forms generated by the rule are
a. valid
b. partially valid
c. invalid
2. There are multiple matching join rules for the pair and they yield the same combined form. The combined form is
a. valid
b. invalid
3. There are multiple matching join rules and they yield the different combined forms. The combined forms are
a. valid
b. partially valid
c. invalid

In order to detect the correct join rule and the correct join forms accurately, two options were considered.

1. Using another set of rules, which can be applied on top of the standard join rules to eliminate false positives. e.g.: When joining the naLu (๑v) and a(\%), the Dvitva Rūpa join rule is also selected. It generates the form nalla. (๓eُeo). But there is an elimination rule that says the letter $L$ (e) is not duplicated. Therefore, the form nalla is eliminated.
2. Check against the set of all valid Sinhala words, so that you can eliminate invalid Sinhala words.

An attempt was made to collect the language rules that can be used to detect the correct join operation for a given pair. But it was found that the documented secondary rule set is not complete, so that in some scenarios, access to the Sinhala vocabulary is needed to check the validity of the combined forms.

Also, an attempt was made to learn these extra rules from a sample data set with tuples of left, right and combined forms. Sinhala being a low resource language, it is difficult to collect an accurately enriched data set large enough to perform the learning to learn the complete rule set.

Having access to a database of all valid Sinhala words is also not practical. Also there are some valid words generated as combined forms by the join rules, that are not valid combined forms of the given 2 words or morphemes.

Hence a combined solution is proposed. It involves finite state transducers for each join rule and non-tagged corpus of Sinhala words.

## 5 Methodology

In this research, we implemented a generic Sinhala word joining tool based on the 9 base rules and 4 additional join rules that are currently in use.


Figure 1 shows the bird's eye view of the word joining process.

Due to the nature of the Sinhala join rules and exceptions in the Sinhala language, finite state
transducers in isolation cannot solve the problem accurately.

For a given left right word/morpheme pair, the joiner applies all the applicable join rules. Some finite state transducers arrive at end states and yield combined forms. All the pairs of combined forms and the join rules used to generate them are returned as intermediate results. There can be both false positives and true positives among them.

A scoring algorithm is introduced to evaluate all the combined forms generated by the join rules. The purpose of the scoring algorithm is to assign a score to each combined form generated by finite state transducers.

The evaluator then selects the results with a score larger than a threshold. The best value for the threshold with respect to a given scoring algorithm is obtained by regressing the joining operation with different threshold values for a sample data set with a manually verified results set.

### 5.1 Scoring Algorithm

The following parameters are passed to the scoring algorithm.

1. Left most word or morpheme
2. Right most word or morpheme
3. Combined form of the left and right.
4. Name of the join rule used to generate the combined form

The algorithm returns an integer value as the score for the given quadruple.

Our software application uses a corpus and some hand coded elimination rules to derive the score.

It first checks whether the combined form is an invalid joined form of the left and right words or morphemes according to the elimination rules. If it is invalid, the score is set as -1 .

If the combined form is not invalid, the word is looked up in the corpus. The occurrence frequency of the word is set as the score. It is a non-negative number.

For the current corpus, the threshold is set as 2.

This value has will depend on the size and quality of the corpus. If this set to 0 , the number of false positives increases due to the impurities in the corpus. If this is set to a larger value, the number of false negatives increases since the evaluator tends to reject valid combined forms that have a low frequency of occurrence in the corpus.

New scoring algorithms may be plugged-in to the application to obtain better results.

## 6 Results

The precision and recall were measured for the joining results of the following data sets.

### 6.1 Dataset 1

8 different grammatical forms of 50 Sinhala nouns were generated by joining their lemma and relevant suffixes. 412 noun forms are expected for the 400 word-morpheme pairs.

### 6.2 Dataset 2

50 pairs of complete words are joined to generate combined words.

### 6.3 Precision and Recall

True positives are the correct combined forms for a given pair generated by the application as the end results.

False positives are the incorrect combined forms for a given pair generated by the application as the end results.

False negatives are the expected combined forms for a given pair, but not given by the application as end results. Some of them were eliminated by the scoring algorithm.

|  | True <br> positives | False <br> positives | False <br> negatives |
| :--- | ---: | ---: | ---: |
| Dataset 1 | 401 | 22 | 9 |
| Dataset 2 | 46 | 0 | 4 |
| Total | 447 | 22 | 13 |

Precision $=$ True positives / (True positives + False positives)

$$
\begin{aligned}
& =447 /(447+22) \\
& =0.9531
\end{aligned}
$$

Recall $=$ True positives $/$ (True positives + False negatives)

$$
\begin{aligned}
& =447 /(447+13) \\
& =0.9717
\end{aligned}
$$

## 7 Conclusion

### 7.1 False Positives

The analysis showed that the main reason for the false positives is the lack of elimination rules.
E.g.

The input : ali $+\bar{a}(\underset{q}{\text { q® }}+\underset{\text { čo }}{ })$
ali is the lemma of the noun elephant
$\bar{a}$ is the suffix to form the singular subject form
Expected output: (aliyā - ̨̨లి $\omega$, āgama)
 ச̧C゚®コ, dwithwạ rūpa)
aliy $\bar{a}$ means elephant. all $\bar{a}$ means god Allah.
The word all $\bar{a}$, though it occurs frequently in the corpus, is not a valid combined form of the lemma ali and the suffix $\bar{a}$.

Some false positives are due to the impurities in the corpus.

We may add an exceptions database and introduce more elimination rules to the scoring algorithm to reduce the false positives.

### 7.2 False Negatives

The analysis showed that the main reason for the false negatives is the incompleteness of the corpus. The occurrence frequencies of the valid combined forms that are not available in the corpus are set as 0 . Therefore, they are eliminated by the evaluator.

It is not possible to create a corpus with all valid Sinhala words. Therefore, we may use statistical or machine learning methods to learn further scoring rules.

### 7.3 Performance

Since our corpus contains 1.2 million entries (including impurities) the database lookup takes a considerable time on test machines. Therefore, an average join operation for a given word pair takes around 20 milliseconds on a 2.5 GHz processor.

### 7.4 Future Enhancements

A possible future enhancement would be to generate a large sample dataset of tuples of left and right words or morphemes, combined forms and rule name using the current word joiner tool version, get them verified using human input and use that dataset to mine the elimination rules using statistical methods.

The elimination rules mined by this exercise may also be used to implement a scoring mechanism for words that are not available in the corpus.

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