## 1. INTRODUCTION.

In this paper I describe a system for the on-line semantic analysis of texts of up to paragraph length. It was programmed and applied in Q32 LISP 1.5 to material of two sorts: newspaper editorials and passages of classical philosophical argument. The immediate purpose of the analysis was to resolve the word-sense ambiguity of the texts: to tag each word of the texts to one and only one of its possible sonses or meanings, and to do so in such a way that anyone could judge the output's success or failure without knowing the coding system. The system tackles texts of up to paragraph length because I take it as a working hypothesis that many word-sense ambiguities cannot be resolved within the bounds of the conventional text sentence; there simply isn't enough context available.

The system attempts to detect semantic forms (which I call templates) directly in coded text, and not by means of a conventional syntax analysis. This restriction sets the present approach apart from the better-lnown ones. However, an approach like the present one still has to show how to obtain the information contained in a conventional syntax analysis, and I shall do that below. For each paragraph of text examined the system dorives a nested structure of the sementic templatos, which can be thought of as its semantic representation. is I shall show, it may be necessary for the system to enlarge its own dictionary in an on-line mode in order to obtain such a representation. From a representation, a word-sense resolution cf the text is read off and printed out, since the representation contains one and only one sense representation for each constituent word of the text.

## 2.

The basic item, the template, is intended to express, in coded form, the message content of an elementary clause or sentence. Thus, if we had to analyse the sentence "The old postman is angry", I would expect to match with it a template that could be interpreted as "A certain kind of man is in a certain state". Similarly, if analysing the clause "The wicked wizard", I would expect to match with it a template that could be interpreted "a man is of a certain kind". The main hypothesis of the system of sense analysis is that one can build up a 'proper semantic sequence' of such templates as a representation of "semantically compatible" fragments of text. At the end of the paper I shall discuss the possibility of explicating the difficult notion of "meaningful language". But at the beginning I am assuming that, if a text is meaningful then its parts must cohere together in some stmuctured way, and that "semantic compatibility" might express that way. This working hypothesis will also mean that the word-senses that can participate in such a proper sequence will be the appropriate ones. By "appropriate senses" I nean simply the dictionary word-senses that a translator of the text would wish to distinguish from the inappropriate ones.

By way of example, I shall consider the senantic compatibilities of the fragments of a peragraph to be found in a"Times" editorial in December 1966. Ls given below it has been fragmented by functions whose operations I shall describe . $\because$ but I shall assume that it is comprchensible as a sequence of twelve items:

```
-1 ((BRITAINS TRANSPORT SYSTEPGS ARE CHANGING)
-2 (aND WITH IT THE TRAVIRLIING PUBLICS HLBITS)
    3 (IT IS THE OLD PERMANENT WAY)
    4 (WHICH ONCE MORE IS EMERGING)
    L5 (AS THLI PAGEMAKERR)
    6 (aIRLINES IfTLLY HAVE BEEN LOSIAG TRiIFIC)
    [7 (TO MODERNIZED RAILWAYS)
    -8 (RAILWIIYS AT LAST ARE BEGINNING)
    [ (TO TAKE SONE CARS
    -10 (OFF THE CONGESTED SYSTEITS TO TUKE THE WEIGHT)
    (11 (IF THE NEW IDEAS ARE FORWIARD PRESSED)
L12 (COMML THE OF. CONMUTRR MOVENENT AND DORMITORY
        .RZ& CONGESTION FLO& P.TTLGN COULD BE CHiJGED))
```

Fig.1. A paragraph in fragment form and it's semantic
compatibilities.

Let's now look at possible semantic compatibilities between fragments of the paragroph (marked with braces in the left hand margin of the figure above).

Fragments $1 \& 2$ are sementically compatible (both essentially assert that a structure is of a certain sort: (1) that a system is changing, (2) that a structure is the public's.) This requires that one takes "to be of a certain sort" in its usual wide logical sense to cover such notions as change and movement:
4 \& 5 are semantically compatible (both essentially assert that something is moving in sone way).
$7 \& 8$ are semantically compatible (both essentially assert that the railways are near to us in time in some way). 2 \& 10 are sementically compatible (both essentially assert that sonething is taking or removing something).
$11 \& 12$ are semantically compatible (both essentially assert that some structure is changing or about to change).

Notice that semantic parallelisms of this sort between fragments are sufficient to resolve at least one ambiguity in each of the pairs of fragments: for example the correct sense of "habits" for fragment 2 is "structure of behaviour", rather than the less-common "articles of dress". Thus pointing out this parallelism is also selecting the appropriate sense of "habits".

## 2. THE TEXTS AND SEMANTIC DICTIONARY

Ten paragraph length texts were chosen for analysis: five from randomly chosen Times"editorials (data texts); and five from the works of philosophers, Descartes, Leibniz, Spinoza, Hume and wittgenstein. The reason for the choice of this type of material will emerge in the discussion. Each paragraph was stored as a list of sentences on a LISP file, and an alphabetical concordance for the texts was obtained with the aid of standard routines. From this the semantic dictionary was written.

The information stored for each dictionary entry word is a list of pairs, each member of which consists of a left-hand member, which is a semantic formula such as (((THIS POINT) TO) SIGN) THING), and a right-hand member, which is a sense description of the meaning of the corresponding formula, such as (COMPASS AS INSTRUMENT POINTING NORTH). Bech such pair (called a sense-pair) corresponds to one sense of the dictionary entry word. The sense description (rigit-hand member of pair) serves only to explain to the operator, in ordinary language print-out, which particular sense of the word is being operated on at any giveñ stage of the procedure. The sense
descriptions are not used as data for computation, except for looking at their first item to get the name of the word in question.

The purpose of the formulae is to encode, and so distinguish the different sense: of natural language words: one would expect to assign a different formula to each major sense of a word that a good dictionary distinguishes. Formulae consists of left and right parentheses and elements, where an element is one of the following 53 primitive semantic clessifiers, or markers;

BE BEAST CAN CAUSE CHANGT COUNT DO DONE FEEL FOLK FOR FORCE FRON GRAIN HavE HOW IN KIND LET LIFE LIKE IINE MLN MAY MORE MUCH MOST ONE PAIR PART PLAMT PLEASE POINT SMME SELF SENSE SIGN SPRELD STUFT THING THINK THIS TO TRUE UP USE WANT WHEN WHERD WHOLE WILI WORLD WRAP.

These elenents constitute the major categories of the classification of word-senses. The whole class of elements is not chosen at rendom; though as with any system of semantic markers it is difficult to justify its memborship in detail on theoretical grounds (though see 4). I shall assume here only that one has to choose some set of markers to work with, and anyone's set of markers is always open to detailed objection. The markers are the basic elements in terms of which all the others in this system (templates, formulae etc.) are defined. So they cannot themselves be further defined, except by means of a table of 'scope notes: which gives the dictionary maker some indication of the marker elements. The table contains entries like:
GREIN: (II,IV,VI) any kind of structure or pattern.
(III) structural or pattern-like.

The Roman numerals refer to the six types of bracket groups used by the dictionary maker in constructing formulae. They are, in order, Ldverbial Group, Adverbial Clause, Adjunctive Group, Nominal Group, Operative Group, Operative Clause. The first two, for exanple, can be illustrated as follows:
I. Adverbial Group.
((TRUE MUCH) HOW)--equivalent for "enough" used as an adverb; same function as "rather nicely" in Fnglish; can end with clement HOW.
II. Adverbial Clause
(NiN FROM)--same function as "out of sight" in English; cannot end with any of the elements of D4 below, and hence a II type cannot be a wellformed formula (see below) by itsclf.

All these six types of sub-parts of formulae can themselves be interpreted (as can the formulac) so that each left-part is dependent on the corresponding right-part. This is a non-intuitive order in LISP but is an aid to reading the formulae for English speakers. This is best explained by means of an example. Thus, to take a sense-pair at random, say (COLOUALESS (( ( ( WHMRE SPREAD) (SENSE SIGN)) NOT HAVE) KIND) (COLOURLISS AS NOT HAVIIG THE PROPEMTY OF COLOUR)))). inn explanation would be; "Colourless" is a sort; a sort indicating that something does not possess some property; the property is an abstract sensuous property of a certain sort; that certain sort has to do with spatial
$\therefore$ :t is not difficult to see that that is what (in rightleft order) the formula conveys.

Formulae are defined recursively as follows:
D.1. A formule is a binarily bracketted string of formulae and atoms.
D.2. An atom is an element, or an element immediately preceded by "NOT".

It follows from this that an element is not a formula. Not all formulae can be assigned to sense-pairs, but only well-formed formulae:
D.3. The head of a formila is its last atom. (and so is the opposite of the usual notion of 'head' in LISP 1.5).
D.4. A well-formed formula (wff) is (a) a formula, and (b) such that its head is one of the following elements: HOW KIND FOLK GAIN MAN PART SIGN STUFF THING WHOLZ WORLD BE CLUSE CHANGE DO FEEL H.VE PLEASE FLIR SENSE WINT USE THIS.
3. INITI. FR FRGMENTATION OF THE TDXTS.

An initicl set of functions breaks each sentence of a paragraph up into strings of words, and, in certain circumstances, reforms discontinuous sub-strings into whole strings. The output from this process is a sentence in the form of a List of "sentence fragments", each of which (if it is not a single word) is either an elementary sentence, a complex noun phrase, or a clause introduced by a marker (such as a preposition).* So for example, the first paragraph of text is returned as on p. 2 above by a function which applies the setfragfentations to each of the sentences of a paragraph in turn, and returns the paragraph as a single list of such substrings, thus obliterating the original sontence boundaries.

* These markers are largely derived from Earl. (3)

It can be seen from the example paragraph above that the functions described do not simply segment sentences in a linear nanner. They also 'take out' certain kinds of clause from within a sentence and append them as separate sub-strings. An example of this 'taking out' and reforming can be seen in the example paragraph reproduced above. The first two fragments read ((BRITiINS TRANSPORT SYSTELS ARE CH:INGING) (ND WITH IT THE TRiVELLING PUBLICS HABITS)).
These are produced from a sentence that oricinally read "Britaine transport system and with it the traveliing publics habits are changing". This sort of break-up leads to an apparent grammatical 'howler', nanely a singular subject for a plural verb. But for the purposes of semantic analysis by the present system that is not a disadventage: it is more than outweighed by having the text cut into serantically acceptable units (see Halliday(4), for the attachment of templates to them.

The fragmented paragraphs are not passed directly to the templatematching procedure, but are first processed by a set of re-ordering functions. These inspect the fragmented out put for a paragraph and seek for qualifying phrases boginning with rarker words like 'of' and 'for'. These are delimited at their other end by the character 'fo', and are placed as a whole before the word they qualify/ adjectives before the preceding noun and so on. Only after this rearrangement are the fragments passed on to the matching functions. The reason for the re-ordering is that when a termplate has been matched with a fracment, the subsequent routines seek for the qualifiers of a noun or verb only to the left of it. Thus a phrase "a book of rules" goes to the matching routines as "a of rules fo book".
the purpose of the fragnent unit is to define a unit of context between the word and the sentence, as usually understood. I shall call "internal" those senantic routines which operate wholly witrin fragmente, and "external" those whiuh scan text outside ${ }^{2}$ particular fragment in order to resolve it.: word-senses.

## 4. THE SYSTEM OF SEMCNTIC HNALYSIS.

## Production of sinsle baro templates

The present system replaces each fragment of text by a number of strings of formulae (frames) oonstructed from the formulae for the words of the fragmens. It then searches yach frame and replaces it by a number of matchine templates, or meaning structures. One can display these procedures schematically as follows:


Fig. 2. ittachment of text to templates.
In the course of these procedures, therefore, each fragment of text is tagged to a number of templates, and so each such template is tagged to some particular selection of the wordsenses for the words of a fragment. line purpose of the subsequent procedures is to redus this "fragment ambiguity" by specifying a set of striags corresponcing to each text fagh, : the so specifying a particular sei. of wort sensed $i x$ worme of the whole text.

The intuitive goal is that there should be just one string of templates in the set, and hence a unique ambiguity resolution of the text. However, the possibility of a number of independent resolutions cannot be excluded a priori.

Thus the outcome of applying these procedures to a text is either nothing, or a string of sense-explanations for the words of the text. In the case where the outcome is nothing, further procedures are defined whereby the system returns, as it were, to the beginning, adjusts one or more dictionary entries in a determinate way and then tries again to resolve the text. Thus the positive outcome described may be achieved after any one of a finite number of tries. As will be seen, there is a limit to the number of possible tries; and after it has been exhausted, the system has to conclude that the text cannot be resolved by this particular method.

The procedures of resolution can be put in the form of a set of phrase-structure rules which produce a nesting of frames of formulae from an initial paragraph symbol $P$. The rules are given in their generative rather then their analytic form, but I give the "lowest-level" rulos first, because they are the ones applied at the first stage of analysis. The presentation will thus end up, rather than start, with highest level rules $P+. .$. , where $P$ is a "paragraph symbol" analogous to the sentence marker, $S$, in conventional grammar.

Following what has been said above:
D.5. A frame for a fragment is a string of formulae such that each word of the fragment that has a (non-null) dictionary entry is represented by one and only one formula, and that formula has the same linear order in the as the
corresponding word in the fragment. Thus the set of all frames consistent with this definition (and with the dictionary entries for the words of some fragment) constitutes an initial representation of a fragment in the system.

We can now define the fundamental notion of template.
D.6. A bare template is any concatenated triple of elements that can be produced by Rules $1-6$ below. (The rules 6. are only a sample).

R1. $T \div \mathbb{N} 1+V+\mathbb{N} 2$
R2. $V \rightarrow B E$
R3. N2 $\rightarrow$ KIND, THIS, GRAIN, THING, SIGN.
R4. N1 $\rightarrow$ GRAIN, THIS, THING, PART, SIGN, MAN, FOLK, SIUFF, WHOLE, WORLD.
R5i. ( $\mathrm{N} 1 \rightarrow$ THIS $)+\ldots+\mathrm{N} 2 \rightarrow$ PLLRT, MAN, FOLK, STUFF, WHOLT, WORLD
ii. $(\mathbb{N} 1 \rightarrow$ THING) $+\ldots+\mathrm{N} 2 \rightarrow$ PART, STUFF, WHOLIT, WORID
iii. $(N 1 \rightarrow P . R T)+\ldots+N 2 \Rightarrow$ PLRT, STUFF, WHOLE, WORID
iv. (N1 $\rightarrow$ SIGN) $+\ldots+\mathrm{N} 2 \rightarrow$ PART, STUFF
v. (N1 $\rightarrow$ MNN) $+\ldots+\mathrm{N} 2 \rightarrow$ FRRT, FOLK, STUFF, MiN
vi. ( $\mathrm{N} 1 \geqslant \mathrm{FOLK}$ ) $+\ldots+\mathrm{N} 2 \rightarrow$ PART, MLN, FOLK, STUFF
vii. $\quad(\mathbb{N} 1 \rightarrow$ STUFF $)+\ldots+N 2 \rightarrow$ PIIRT, STUFF, WHOLE, WORLD
viii. (N1 $\rightarrow$ WHOLE) $+\ldots+N 2 \rightarrow$ PART, STUFF, WHOLE, NORID
ix. $\quad(\mathbb{N} 1 \rightarrow$ WORLD $)+\ldots+\mathbb{N} 2 \rightarrow$ P.RT, STUFFF, WHOLE, WORLD
x. $\quad(\mathbb{N} 1 \rightarrow G R 1 N)+\ldots+\mathbb{N} 2 \rightarrow P_{i L R}$

R6i. ( $N 1 \rightarrow$ GRiIN $)+\ldots+V \rightarrow P_{4} I R, ~ D O, C . U S E, C H . N G I, H: V E$
ii. $\quad(\mathbb{N} 1 \rightarrow$ THIS $)+\ldots+V \geqslant P A I R, D O, C . J T E T, ~ C H . N G E, H / V E$

The form of rules 5 and 6 is simply a onvenicat abbreviation of a more conventional form. For exampl.

R5 iv. (N1 $\rightarrow$ SIGN) $+\ldots+N 2 \rightarrow$ Pirm STUFF

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is simply an abbreviated expression of the two context-
dependent phrase-structure mules:
        SIGN+\ldots+ N2 }->\mathrm{ SIGN+...+ PIRT, and
        SIGN+\ldots+ N2 }->\mathrm{ SIGN+...+ SIUFF.
        These rules produce bare templates in the form:
        Substantive (or noun) type element +
        Active (or verb) type element +
        Substantive (or noun) type element.
        Thus MiN+H.vE+PNRT con be produced in this way, but
1LN+BE+WORID cannot. This order we call the standard order,
and templates are always considered and compared in this
order even if located in fragments in other (nonstandard)
orders, or in "debilitated forms."
D.7 & 8. If N1+VN2 represents the standard order, then
        V+N1+V2 and N1+N2+V are nonstandard orders, and
        N1+N2
        N1%.V
    N1 N1
        V are debilitated forms.
D.9. A fragment matches with templates if a frame for it
contains concatenations of heads (in left-right order)
corresponding to any template produced by Rules 1-11.
Where: (* indicates a blank litem).
R7: THIS >**
R8: B2 }->
R9: KIND }->\mathrm{ *
R10: V V *
```

```
R11.i N1 +\ldots..+(KIND }->*)->\mathrm{ KIND +...+ N1
    ii. (V >*) +\ldots+ IMIND }->\mathrm{ KIND +...+V
    iii. N1 +\ldots..+(V ( }->*\mathrm{ *) }->\textrm{V}+\ldots+N
    iv. (v 谟)+\ldots+N2 }>\textrm{N}2.+\ldots
```

Rules 1-6 produce standard forms of bare template, and Rules
7-11 praduce (by means of deletions and reordiringi) the
permitted debilitated and nonstandard forms. The latter
rules produce actual text-items, in the sense of heads (of
formulae) to be located in the frames that represent fraements
of text directly.


Fig 3. Preference table for bere tempiates.

Since Rules 1-11 are nonrecursive, there is no problem about orderine the productions in this way. Apart from the forms given in the table, there are only vacuous cases such as *+*+*.

The above table is intended to make clear the relation between the various standard forms (in the rightmost column) and the corresponding "items in frames" produced or recognized (middle column). Thus in the generative mode, text items are produced from the standard forms by transposition and deletion. In the analytic mode the text-items are recognized in the rank order shown, and then transposed and augmented with dummy $B P$ and THIS elements so as to be in standard form for further computation.

The actual function of the rank choice is best explained by example, particularly as regards the composition of Rank I, since the ranks lower than I clearly consist of "debilitated forms" and it is intuitively plausible to produce fuller forms first. This ordering is one example of the general rule which enables template matching to do (at least) the work of a conventional grammar; namely, pack the frome. as tightly as possible, or, in other words, produce the fullest possible template.

The presence in kank I of the debilitated form KIND $+\mathbb{N} 1$ can be understood by considering, for example, the fragment:
(THE OLD TRAISPORT SYSTEN).
To simplify matters I shall consider only (i) the frame consisting of representations of the appropriate senses of the words in that fragment, and (ii) the frame identical with the first except that it contains representations of OLD as substative (noun $=$ "the old people") and the active (verb)
form of TRANSPORT. Thus, by the semantic coding system described above, those two : will contain the following heads, and in the order shown:

```
i. ... ...KIND) ...KIND) ...GRAIN) , and
ii. .. ...FOLK .....DO) ...GRiIN) .
```

Now the above rules generate both
(FOLK+DO+GRAIN) and (KIND+GRAIN)
as strings of text-items; the latter by deletion from ( $\mathrm{N} 1+\mathrm{BE}+\mathrm{KIND}$ ) and (KIND +N 1 ). It is clear that if the form KIIJD $+\mathbb{N} 1$ were not in Rank I with forms like ( $N 1+\mathrm{V}+\mathrm{N} 2$ ) which yield (FOLK+DO+GRAIN), then a substantive phrase like this one would never receive a proper interpretation, since Rank I (without the form (KIND+N1)) would always look for an active (verb) sense for"transport"and having found one, would be satisfied.

As I have described the process so far both bare template forms (FOLK+DO+GKAIN) and (GKINT+BE+IIND) would be produced. I shall show in the next section the additional procedures which produce the second of these in preference to the first Production of single full templates.

Further production rules limit the templates actually produced, and these require the notion of full template, defined as follows:
D.10. $A$ full template is two triples of formulae such that the heads of the first triple constitute a bare template, and the second triple can be produced from the first by the rules 12-16.
D.11. The six formulae constituting a full template are called text-values.

The six formulae so defined give content to the corresponding bare template (expressed by the heads of three of the formulae). The rules $12-16$ specify the other three formulae in such a way that each of them can be the qualifier of one of the formulae with a head defining part of the bare template. The rules 12-16 (not given here for reasons of space) are, in effect, rules producing an ordered pair of formulae such that the first is an appropriate qualifier for the second. Thus rule $13 i$ produces an adjective type of formula (one ending in KIND) before a noun-type of formula, and so on.

The full templates are the items with which the system really operates. They can be illustrated by contrast with bare templates by considering fragment 3 of the paragraph excmined earlier. That fragment was "It is the old permanent way". Amons the bare templates produced for it by the system are the following two:

```
((IT IS THE OLD PERMGNENT WaY)
    ((THIHG BD SIGN)
        (((THIS THING) (IT AS INANINLTE FRONOUN))
                ((BE BE) (IS AS HLS THE PROPBRTY))
                ((((NAN FOR) ((WHWHE POIITT) FROMI)) (LINE SIGN})
                (WAY IS PaTE OR ROUTE))))
((THING BIL SIGN)
    (((THIS THING) (IT AS ININIM.TE PRONOUN))
        ((BE BE) (IS .S HAS THD PROPERTY))
        ((((THIS THING) (TRUE USE)) SIGN) (NGY NS MEINSS))))
```

The fragment here is tied to two items, each of which is a bare template triple followed by the three formulae in the sense frame which locate it (their last elements are the same as thoss of the template triple point of

## 17.

interpretation should be added here for speakers of .mericen English: all speakers of British Bnglish interpret "way" in this fragment as heving its "path or route" sense in trise context.

The two bare templates are now expanded to full templates as follows:
((IT IS THE OLD PERRLINEMT WAY) ((THING BE SIGN)
(( (ONE THING) (IT LS INLMINaTE PRONOUN) )
( (BE BE) (IS LS HLS THE PROPERTY))
((( (THIS MHING) (TIUE USTE)) SIGN) (W.Y IS ME:NSS))
NIL NIL ((NOTCH_NGE KIND) (PERMMNENT LS UNCELNGING))))
((THING BE SIGN)
(( (ONF THIING) (IT AS ININNMMTE PRONOUN))
( ( BE BE ) (IS I S HLS THE PROPRRTY) )
((( WHERE IN) ((WHERE POINT) FKOM )) (IINE SIGN))
(WiY IS PiATH OK ROUTE))
NII NIL ( (NOTCHiNGR KIND) (PGRMNENT AS UNCH.NGING)))))

These two items are the expansions (in franes of sense pairs) of the two bare templates. They consist of the same items as the bare template plus three formulae which are the qualifiers of the first three, (the fourth of the six is the qualifier of the first of the six and so on). In this the 'it' and 'is' have no qualifiers, honoe the LISP 'INII's in those positions. Bare templates other than these two were matched onto the fragment, but only these two could be expanded in this way. Hence these two were the 'survivors' and the others were rejected from further consideration.

Wen expanding in this way to produce full templates from bare ones the following metamule (:15) is applied

```
"Produce preferentially those full templates in which as
many elements as possible are developed by the rules R12-
R14!' This means producing if possible those full templates
in which each element of the bare template has a formally
appropriate predecessor. By means of a further rule (R16)
an attempt is made to produce not only full terplates with
formally appropriate internal relations, but also ones with
semantically close internal relations as well. That is to
say, full templates such that the triple of qualifying
formulae are senantically close to the formulae they respect-
ively precede. Where,
D.12. Two formulae are said to be semantically close if:
i) they share a common pair of elerients; or
ii) they have one or more of the following elements in common: ONE, COUNT, WORLD, WHOLE, LIFE, IIND, NTUST, SELF, SPREAD, TRUE, WR.P, WHEN, WHAFE, THINK; or
iii) Their cores are such that they are identical, or either is a nember of the other in the sense of a list-nember, or the left or right hand member of either core is a member of the other.
```


## Rules producing more than one template

I can now consider the production of concatenations of the full templates described so far.
D.13. A paragraph string is any string of templates produced by the rules $17 \& 18$ from the paragraph symbol $P$.

R17. $\quad \mathrm{P} \rightarrow \mathrm{T}_{\mathrm{r}}+\mathrm{T}_{\mathrm{s}}$
if $T_{r}$ is a full template written as a string of six formulae thus,
19.

$$
\begin{aligned}
& \left.F_{r 1}^{1}+F_{r 1}+F_{r 2}^{1}+F_{r 2}+F_{r 3}^{1}+F_{r 3}\right] \\
& \text { where } F_{r 1} \text { is a noun type; } F_{r 1}^{1} \text { its qualifier (adjective type); } \\
& F_{r 2} \text { is a verb type, } \mathrm{F}_{\mathrm{r} 2}^{1} \text { its qualifier (adverb type), and so } \\
& \text { on, then } \\
& \text { R18. } \\
& \left(T_{s} \rightarrow F_{s 1}^{1}+F_{s 1}+F_{s 2}^{1}+F_{s 2}+F_{s 3}^{1}+F_{s 3}\right) T_{s} \\
& \rightarrow\left(F_{t 1}^{1}+F_{t 1}+F_{t 2}^{1}+F_{t 2}+F_{t 3}^{1}+F_{t 3}\right)+\ldots+ \\
& \left(F_{u 1}^{1}+F_{u 1}+F_{u 2}^{1}+F_{u 2}+F_{u 3}^{1}+F_{u 3}\right),
\end{aligned}
$$

where the values of the two template forms produced are semantically close.
D.14. Two full templates $T_{r} \quad T_{S}$ are semanticelly close if (with the above notation for full templates) at least two of the following pairs of formulae are (i) such that the head of the second is identical with, or in the neeation"class of, the first:
$\left.\left(F_{r 1} F_{s 1}\right),\left(F_{r 1} F_{s 3}\right),\left(F_{r 2} F_{s 2}\right),\left(F_{r 3} F_{s 1}\right), F_{r 3} F_{s 3}\right)$; and (ii) either they, or their qualifier formulae, are senantically close. These ten possible directions of connection between two full templates can be shown schenatically as follows:
$\left.F_{r 1}^{1}+F_{r 1}+F_{r 2}^{1}+F_{r 2}+F_{r 3}^{1}+F_{r 3}\right]$ qualifier $\mathbb{N} T Y P E$ qualifier $V$ TYPE qualifier N TYPE


Fig 4. Connecting pattern betwe, full templates.
see note on page 31.

Rule 18 does not, as might appear at first sight, involve self-contradiction. The sharthand form of rule writing is now being extended to mean that when $T_{S}$ has been rewritton as $F_{s 1}+\ldots \ldots .+F_{s 3}$, then the latter may be rewritten as the right-hand side of the second arrow.

This "expansion-concatenation" rule can be recursively applied to the initial productions from $F$. Thus at any stage in the process a paragraph string of full templates is produced. It any point the string can be considered terminal and, with the aid of the dictionary of words and sense-pairs, the paragraph string of templetes can be converted to a string of frames and so to a text of words. This is analagous to the introduction of the lexicon in any standard phrasestructure cramnar. The dictionary entries themselves can be put in phrase structure form. For example, if a word $W_{n}$ has two sense pairs $S 1$ and $S 2$ in its dictioncry entry, then the sense-pairs thomselves cen be put in the form $S 1 \rightarrow W_{n}$ and $S 2 \rightarrow W_{n}$ respectively. This form of the dictionary entries is useful in representing the self-modification of the system described below.
5. APPLICATION OF THE SYSTEM TO TEXPS.

## Matching bare templates onto fragments.

Rules 1-6 above define the matching of bare templates onto a fragmented text, one bare template onto each text fragment. TEMPO is the main (top level) function that does this: it examines in turn all the frames of sense pairs for a fragment, and so on for all the fragments of a paragraph.

It takes as its argument a frame of sense-pairs, one for tach word of a given fragment. THMPO scans each such combination in turm, starting with the frame containing ell the main senses of the words in the fragment (the first ones in the dictionary entry for each word). TMMP searches for triplets of heads in the order of preference given in fig. 3 , above. For example, if it finds type $I$ templates it doesn't look for any of types II-IV and so on. Each type of template is collected on a list which is the value of a different free LISP variable. If TEMPO finds nothing till it reaches the debilitated $\mathbb{N}+\mathbb{N}$ form, it replaces the $\mathbb{N}+\mathbb{N}$ by $\mathbb{N}+B E+N$ ( $B E$ being the "dumy verb"). Similarly $V+\mathbb{N}$ and $\mathbb{N}+V$ are replaced by THIS $+V+N$ and $N+V+T H I S$ respectively (THIS keine the "durmy substantive"). The function of these duram features is to supply a general form of template for subsequent processing, even when it is not wholly present in the text. Suppose, for example, a fragraent consisted not of an assertion form, but of a noun phrase like "the black wizard", whore the heads of the appropriate codings for "black" and "wizard" would b.. KIND and MiN respectively. As there is no verb, a debilitated template of the $N+\mathbb{N}$ form would match onto these two heads, and that would then be converted into MNF+BE+KIND. which is the intuitively correct interpretation (WIZARD is BLiCK). The dummy verb is added in the way described; and in cases like this, where the first head is the predicate KIND, the order of the two heads is reversed, so as to give the $M N+B E+$ KIND form. This transposition is defined by R11i.

The internal rejection functions (matching full temrlates)

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Earlier I distinguished between intornal and externel procedures. Internal rejections are those procedures which cast out matchine templates ly means of the
expansion from bare to full templates. The main function which does this is PICKUP. It takes a fragment name as argument and constructs the TEMPO value for it. PICKUP makes a decision in the case of each templets whether or not to reject it from further consideration. Those that survive are then considered further by the externel rejection procedures. The survivors from PICKUP represent a stage of ambiguity resolution beyond that given by TEMPO. If, for example, PICKUP examines a template that has been natched onto a fragment containing the words round box, where a template head had been attached to a formule for box', then, hopefully, PICHUP keeps at least the terplatc in which round is codod by its "spatial property sense" and box" is coded by its "container" sense.

Inside PICKUP the function REFINE returns as its value a list of five sub-lists of full templates: its first sublist contains those form-close internally in four ways (as defined by rules 12-15), down to the last sub-list containing twose with no such closeness. PICKUP takes the first non-empty submist of REFINE, and of that returns as its value the full templates that are semantically close as well (if any).

The 'semantic parser'; resolving a paragraph.

The top-level function PiRSPARA takes as its argument a list of fragments, produces the PICKUP value for each (in the full template form given on $p .14$ ) and then parses these full templates using rules 17 \& 18 . A nesting of templates that satisfies these rules is an interpretation for the paragraph, and its word-sense content is read off and printed out (since a nesting of full templates is simply a selection of the possible word-sense assignments for a text.) Full templates which cannot be parsed with those for other
fragments are simply rejected . This is the aternal rejection procedure referred to cerlier.

Functions called FIT and JLM express rule 18: they test for semantic closeness between two full templates and, if such closeness is found, the two full templates are replaced by a single item with the form of a full template. Or - to put it in terms of the two function nomes - if the full templates FIT, they are then JAMmed. If the three main formulae in a full template are related to the three main formulae of another template by any three of the connectivities expressed in fig.4. above, then the two templates FIT (are semantically close). The function JIM builds up a representation of the two templates based on their connectivities. FIT and JAM work with pessege-pairs, which are to a fragment what a sense pair is to a word.
D.15. A messcoge-pair is a two-item list: one item is a list of the first three sonse-pairs of some full template, the other item is a list containing the name of some fragment with which the full template metches.

PirSPARh constructs the PICSUP value (full templates) for its list of fragments, and then builds up all possible frames of message-pairs for the paragraph. Each frame of message-pairs is now a possible meaning representation for the whole paragraph. PARSPARA then scans each frame in turn to see if it can find a right-left contiguous pair of message -pairs satisfying FIT. If it cen it daletes the first message-peir and replaces the acond by a messaze-pair consisting of (1) the JMM value of the tro 'parsed' full templates, and (2) a list of the names of he fitting fragments.
(pusitg .à Gunecmid mith the whols feofia))
NIJ (( (wBELA CH:NGB) IOW)

then the two full templates in those mosseftepairs are a fittine pair, it whall expoct then to be replanod in the strine by the form:

(with in the raveluing fublicis hibins))

((BE BE) (ific is HaVe pre proplity))
((Chiagt kind) (changing is inmziliag))
(( (Thing for) ((wricio Clinge) kind))
 NIL (( $($ Whaze Ch.MGE) HOW)


This fittine together, or parsine, of message-pairs axpresses the semantic compatibility between the corresponding fragments discussed earlier. PiESPaku rewrites such strings
of message-pairs recursively, trying to reech a two item list which (by rule 17) is $P$ the paragraph symbol. If this point is reached the corresponding sense-resolution is read off and printed out for the paragraph in the following form: each fragment is given with the list of sense expressions for all the words in it which are resolved (or which had only a single sense entry initially, and so are trivially resolved); a list is also given of words not resolved (if any).

((WORDS RESOLVED IN FRiGMEMT)
((TRINSPORT AS PERTANING TO MOVING THINGS :BOUT) (BRITAILS AS HAVING THE OHNELCTERISTIC OF A PLITICUI:AR PiAT OF THE WOLDD (SYSTEN IS IN OKGANIZSTION)

((WORDS NOT RESOLVBD IN FRAGNENT) NIL))
( (WITH II IMI: TAVELLING PUBLICS HABITS)
((WORDS IEESOLVED IN FRIGMEMT)
( (TRAVELLIING aS MOVING FROM PL:CE TO PLACE)


((WORDS NON RBSOLVED IN FRLGMENT) NIL)
fig. 5. First two fragments of the resolved output for e text paragraph.

The orieinal English for the first two fragments of that paragraph was "Britain's transport system and with it the travelling public's habits are changing".

## The sense constructor procedure.

A procedure was built in to the system to deal with the
cases where the system returned (NO RESOLUTION ALI PATHS BLOCisED) at the teletype. This situation could arise for a number of reasons; the text fragments did not cohere together sufficiently; a vital word sense had been left out of the dictionary; or a word in the text was beinc used in a new and original sense. An obvious suggestion for tackling this is to allow the word dictionary to enlarge itself: to supply an additional sense entry for the word that is holding the procedure up, if it can be found. Such a construction could thought of as adding a new rule F - a, where F is a formula and a word name, and so expanding to a new rule system as the system adjusts to the particular text.

In practice PARSPARA examined the value of a free variable BESTPARS each time it failed to parse a frame completely. It stored as the value of BESTPARS the parsing tree containing the templata that had been rewritten least. It seemed a good first guess at the recalcitrant word that it was in template that 'cohered' least with its neigbours. If all the frame blocked PARSPARA would print (CONSTRUCTER MODE) and evaluate a function of no variakles called CONSTRUCTER. This function controls all subsequent operations via the READ and PRINT functions ai the teletype. CONTSTRUCTER looks ar the value of the recalcitrant template in BESIPAFS and suggested that a word in the corresponding fragment have its dictionary of sense pairs enlarged by identifying the recalcitrant word with the nost 'semantically close' word in the paragraph. If the operator accepts the system's suggestion at the teletype, the system is rerun with the enlarged dictionary to try and get a resolution. In such a case (or if none of the system's suggestions are acceptable to the operator) the system returns to the normal operating mode. This procedure was not called upon for the newspaper paragraphs, but it produced some interesting suggestions in the case of two of the philosophical paragraphs.

In CONSTRUCT: MODE dialogues like the following are possible:
(COnsmructir mode)
((no resolution all paths blockid)
(best parsing contalns)
(()((kind sign) (attribute as a particular kind of PROFERTY))
((BI BIE) (DUNMY))

((MHOLE (MUST (KIND SIGN)))
(Nature as issence of mscivilai propertibs)) NIL NIL)

(ricaicimant mbipleme is foir)
(the shie mature or aithibute)
(continue yes or no)
YES
(SUGGist antiklbute as mature (shail I mry it yes or no)) YES
(( (If Thedia gere tho ca moat disminct substances)
( (words resolviad in Frigminis)
((There As AT A POINT) (WERL LS EXISTED) (OR is DISJUNCTION)
 (distingt as drferilify) (substances as sorts of thing ))
((words not iasolvid in Frighint) (Tw (( (COUNT SIGN) (TWO iS is NUMBEMR)) ((count kind) (two is hiving tee proptrpy of twoity)))))
fig. 6. Dialogue in CONSTRUCTER MODE together with first part of subsequent resoiution.

## 6 DISCUSSION

One of the main difficulties in coding for, and evaluating, a system like this one is the necessary vagueness of some of the sense-entries (especially evident in words like 'it' and 'is'). Noktreless I claim. that the present system could constitute a tentative criterion for meaninyfulness: a text is meaningful if and only if a system like tins present one can resolve it. It is easy enough to get a necessary criterion ${ }^{*}$ on the ground that one needs to be able to tell in what senses the words of a text are being used in order to call it meaningful. I have argued at length elsewhere that it is possible also to justify the corresponding sufficient one (8). The establishment of such a criterion would be of some interest in the cases of the five philosophical paragraphs, since it was texts like these that Carnes: (2), and the 'Logical Syntax' school eenerally, said could be shown to be meaningless on the basis of a system of analytic rules, though they never in fact constructed such a system. The criterion sugesested here would only be one of degree (in terms of the number of applications of the sense-constructer procedure a text required for resolution). That is perhaps the only acceptable form that a criterion of meaningfulness could take, as there seems something absurd about an attempt to set an absolute bound to the meaningful.

Another speculative interest of the present system might be its application to the speech patterns of schizophrenics. Schizophrenic discourse seems (6) to be meaningful within the bounderies of units of the same order of length as the clause or phrase. The troubie is that these units don't seem to fit together in a coherent way in the schizophrenic's
speeck pattern. A system of the present sort, which tries to make such items cohere, might conceivably provide a measure of "sementic disorder" in such cases.
4. number of connexions can be made also between the semantic structure assigned to a text by the present system and that assigned by formal logic. These connexions have been investigatied in the cases of the five philosophical paragraphs, which have a fom sufficiently like the one required by formal logic. These connexions are of some interest in view of the almost total neglect of the sense-ambiguity of natural lancuace words by formal logic.

One can, for example, interpret the present system so as to create a notion of "valid and usefuli" axgument. It has long been recognised that an argument can be forraally valid (and even have true premisses) and yet be conpletely useless. This is usually due to a genuine ambiguity in the argument. For example, the following is perfectly valid: "Lll kings wear crowns, all cnowns are coins, therefore all kings wear coins". And, within the coiltext of each premiss, each premiss is true. (In the "numismatic world of discourse", for example, the second is true).
in argument could be deemed "valid and useful" if it is formally valid and if the present system assigns to it a consistent and complete interpretation。 I an using the terms 'consistent' and 'complete' in a way similar to Bobrow's (1) use of them: an interpretation is complete if the system assigns an interpretation to each key term in the arsument, and consisten if it assigns the same interpretation(wordsense) to every occurrence of a term. Thus the argument above
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would not pass the 'usefulness' criterion, since a proper ambiguity-reaolver would assign dj.fferent interpretations to the two cocurrences of the key term 'crown'.

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te page 19:
The negation class of elements for each element is derived jnductively by a separate procedure. The notion onvclved is like that of logical contrarysan element and any member of its ngeation class are partly synonymous and partly exclusive. For example, an entity can be basically a sTUFY or basically a THING;it cannot be both so each of these clements is in the nagation class of the other.
